# The Virtual Element Method for Polygons with Curved Edges 

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## The General Field Equations



Volcano Eruption


Black Holes


Cell Biology

$$
i \hbar \frac{\partial}{\partial t} \boldsymbol{\Psi}(\boldsymbol{r}, t)=\left[\frac{-\hbar^{2}}{2 m} \nabla^{2}+\boldsymbol{V}(\boldsymbol{r}, t)\right] \Psi(\boldsymbol{r}, t), \quad\left[D_{0}+D_{1} v\right] J \boldsymbol{C}^{-1}+\zeta(\nabla \cdot \boldsymbol{u}) \boldsymbol{I}=\boldsymbol{D}(v, \boldsymbol{F}, \boldsymbol{\Pi})
$$

$$
\left.\boldsymbol{D}(v, \boldsymbol{F}, \boldsymbol{\Pi})=\left[D_{0}+D_{1} v\right] J \boldsymbol{C}^{-1}+D_{0} / 2 J f_{0} \oplus f_{0}+D_{2} F+\sum_{j=1}^{\infty}\left[\sigma_{j}(t, x) \cdot \nabla\right) u(t)+\nabla \tilde{p}_{j}(t, x)\right] d W_{t}^{j} \boldsymbol{F}^{-1}
$$

$$
\mathcal{L}(\varphi)=-\frac{(\partial \varphi)^{2}}{2}-\boldsymbol{m}^{2} \boldsymbol{M}^{2}\left[\left(1+\varphi^{2} / \boldsymbol{M}^{2}\right)^{1 / 2}-1\right], \quad \boldsymbol{\sigma}=\zeta(\nabla \cdot \boldsymbol{u}) \boldsymbol{I}+\mu\left(\nabla \boldsymbol{u}+(\nabla \boldsymbol{u})^{T}-\frac{2}{3}(\nabla \cdot \boldsymbol{u}) \boldsymbol{I}\right)
$$

$$
d \boldsymbol{u}(t, x)=[\nu \Delta \boldsymbol{u}(t, x)-(\boldsymbol{u}(t, x) \cdot \nabla) \boldsymbol{u}(t, x)-\nabla p(t, x)] d t-\left[g\left(|\boldsymbol{u}(t, x)|^{2}\right) \boldsymbol{u}(t)-f(x, \boldsymbol{u}(t, x))\right] d t
$$

$$
R_{\mu \nu}-\frac{1}{2} R g_{\mu \nu}+\Lambda g_{\mu \nu}=\frac{8 \pi G}{c^{4}} T_{\mu \nu}, \quad \rho\left(\frac{\partial \boldsymbol{u}}{\partial t}+\boldsymbol{u} \cdot \nabla \boldsymbol{u}\right)=-\nabla p+\mu \nabla^{2} \boldsymbol{u}+\frac{1}{3} \mu \nabla(\nabla \cdot \boldsymbol{u})+\rho \boldsymbol{g}
$$

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$$
\left\{\begin{aligned}
-\Delta u=f & \text { in } \Omega \subset \mathbb{R}^{2} \\
u=0 & \text { on } \partial \Omega
\end{aligned}\right.
$$

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$$

Joint work with L. Beirão da Veiga, F. Brezzi, D. Marini and G. Vacca

## Poisson equation

Let $\Omega$ be a reasonable domain in $\mathbb{R}^{2}$, and $f \in L^{2}(\Omega)$. We define:

$$
V:=H_{0}^{1}(\Omega), \quad a(u, v):=\int_{\Omega} \nabla u \cdot \nabla v \mathrm{~d} x, \quad F(v):=\int_{\Omega} f v \mathrm{~d} x
$$

And we consider the following variational problem:

$$
\left\{\begin{array}{l}
\text { find } u \in V \text { such that }  \tag{1}\\
a(u, v)=F(v) \text { for all } v \in V(\Omega) .
\end{array}\right.
$$

This is the weak form of Poisson equation with homogeneous Dirichlet boundary conditions:

$$
\left\{\begin{aligned}
-\Delta u=f & \text { in } \Omega \\
u=0 & \text { on } \partial \Omega .
\end{aligned}\right.
$$

We will consider conformal Galerkin approximations of (1), with $V_{h} \subset V$.

## Key steps of the Virtual Element Method

- Decompose the domain into general polygons.
- Define local spaces using differential operators. Ensure that polynomials are inside. Choose carefully the degrees of freedom.
- Define projection operators onto polynomials which are computable from the degrees of freedom.
- Use the projection operators to discretize the problem, ensuring $k$-consistency ( $\approx$ patch test). Add stability.


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## The local VEM space $V_{k}(P)$

Let $P$ be a polygon with straight edges. Mimicking the classical Lagrange Finite Elements for triangles, we define a finite element space $V_{k}(P)$ on $P$ such that:

- $V_{k}(P)$ contains the space $\mathbb{P}_{k}(P)$ of polynomials of degree less than or equal to $k$ (plus - possibly - other non-polynomial functions);
- a function in $V_{k}(P)$ restricted to an edge $e$ is in $\mathbb{P}_{k}(e)$, so that if two polygons $P$ and $P^{\prime}$ have an edge in common, the two spaces $V_{k}(P)$ and $V_{k}\left(P^{\prime}\right)$ can be easily "glued" in $C^{0}\left(P \cup P^{\prime}\right)$ (compatibility with FEM);
- given functions $u_{h}, v_{h} \in V_{k}(P)$, I can "compute" the bilinear form $a_{P}\left(u_{h}, v_{h}\right)$ and the load term $F_{P}\left(v_{h}\right)$ (or at least reasonably good approximations).


## The local VEM space $V_{k}(P)$

We start with a trivial remark:

$$
p_{k} \in \mathbb{P}_{k}(P) \Longrightarrow \Delta p_{k} \in \mathbb{P}_{k-2}(P)
$$

Since we want that our space contains the polynomials of degree $k$, we define the local VEM space as follows:

$$
V_{k}(P):=\left\{v_{h} \in C^{0}(\bar{P}) \text { such that: } v_{h \mid e} \in \mathbb{P}_{k}(e), \Delta v_{h} \in \mathbb{P}_{k-2}(P)\right\}
$$

If the polygon $E$ has $N_{e}$ edges (and $N_{V}=N_{e}$ vertices), it is clear that

$$
\operatorname{dim} V_{k}(P)=N_{V}+(k-1) N_{e}+\operatorname{dim} \mathbb{P}_{k-2}(P)=k N_{e}+\frac{k(k-1)}{2}
$$

Is the space $V_{k}(P)$ OK?

## The local VEM space $V_{k}(P)$

## Is the space $V_{k}(P)$ OK?

YES: it contains polynomials of degree up to $k$ and is conformal, so (under some regularity assumptions on the polygonal mesh) order $k$ approximation in $H^{1}$ is guaranteed. . .

BUT: the functions in $V_{k}(P)$ are not known explicitly; they are defined through the solution of a PDE in the element.

## We need computable projectors!

## Degrees of freedom in $V_{k}(P)$

Recall that

$$
\operatorname{dim} V_{k}(P)=N_{V}+(k-1) N_{e}+\operatorname{dim} \mathbb{P}_{k-2}(P)
$$

As degrees of freedom in $V_{k}(P)$, we choose:

- the value of $v_{h}$ at the vertices and at $k-1$ points on each edge;
- the (scaled) moments $\frac{1}{|P|} \int_{P} v_{h} m_{\boldsymbol{\alpha}} \mathrm{d} x$ for $0 \leq|\boldsymbol{\alpha}| \leq k-2$
where $m_{\boldsymbol{\alpha}}$ is a scaled monomial of degree $|\boldsymbol{\alpha}|$.
We easily see that the degrees of freedom above are unisolvent in $V_{k}(P)$.


## The $\Pi_{k}^{\nabla}$ projector onto $\mathbb{P}_{k}(P)$

We can define a projector $\Pi_{k}^{\nabla}: V_{k}(P) \rightarrow \mathbb{P}_{k}(P)$ which is computable.
The projector $\Pi_{k}^{\nabla}$ is orthogonal with respect to the energy scalar product in $H^{1}(P)$ :

$$
\begin{gathered}
\int_{P} \nabla\left[\Pi_{k}^{\nabla} v_{h}\right] \cdot \nabla p_{k} \mathrm{~d} x=\int_{P} \nabla v_{h} \cdot \nabla p_{k} \mathrm{~d} x \quad \text { for all } p_{k} \in \mathbb{P}_{k}(P) \\
\int_{P} \Pi_{k}^{\nabla} v_{h} \mathrm{~d} x=\int_{P} v_{h} \mathrm{~d} x \quad \text { to fix the constant function }
\end{gathered}
$$

By integration by parts we easily see that $\Pi_{k}^{\nabla} v_{h}$ is computable from the degrees of freedom of $v_{h}$.

## The $\Pi_{k}^{\nabla}$ projector onto $\mathbb{P}_{k}(P)$

We basically only need to check that the right hand side

$$
\int_{P} \nabla v_{h} \cdot \nabla p_{k} \mathrm{~d} x
$$

is computable from the degrees of freedom of $v_{h}$. Integrating by parts:

$$
\int_{P} \nabla v_{h} \cdot \nabla p_{k} \mathrm{~d} x=-\int_{P} v_{h} \Delta p_{k} \mathrm{~d} x+\int_{\partial P} v_{h}\left(\nabla p_{k} \cdot \boldsymbol{n}\right) \mathrm{d} s
$$

and the result follows by observing that

- $\Delta p_{k} \in \mathbb{P}_{k-2}(P)$ which implies that $\int_{P} v_{h} \Delta p_{k} \mathrm{~d} x$ is computable,
- $v_{h}$ is a known polynomial of degree $k$ on each edge,
- $\nabla p_{k}$ is a known polynomial of degree $k-1$ on each edge.


## Polygons with curved edges

Assume that we have a polygon with a curved edge $\gamma$ :


In the rest of the talk:

- $e$ will denote a straight edge;
- $\gamma$ a possibly curved edge.


## The local VEM space $V_{k}(P)$

We make a crucial remark:

- In order to define the local VEM space $V_{k}(P)$, we only need to change the definition of the space on the curved edge $\gamma$.

The rest of the construction can be done in the same way as before.
A function $v_{h} \in V_{k}(P)$ is then defined by the following properties:

- for each straight edge $e, v_{h \mid e} \in \mathbb{P}_{k}(e)$
- for each curved edge $\gamma, v_{h \mid e} \in V_{k}(\gamma)$ (to be defined!)
- $v_{h \mid \partial P}$ continuous
- $\Delta v_{h} \in \mathbb{P}_{k-2}(P)$ as before.


## Does $V_{k}(P)$ still contain polynomials of degree $k$ ?

Does $V_{k}(P)$ contain polynomials of degree $k$ ?
The answer is:
it depends on whether $V_{k}(\gamma)$ contains the restrictions to $\gamma$ of polynomials in two variables.

We define

$$
P_{k}(\gamma)=\{\text { restrictions to } \gamma \text { of polynomials in two variables }\}
$$

If $P_{k}(\gamma)$ is "almost contained" in $V_{k}(\gamma)$
then
$\mathbb{P}_{k}(P)$ is "almost contained" in $V_{k}(P)$.

## Does $V_{k}(P)$ still contain polynomials of degree $k$ ?

Let $p_{k}(x, y)$ be a polynomial of degree $k$ in two variables. We want to see to what extent " $p_{k} \in V_{k}(P)$ ".

Assume that $V_{k}(\gamma)$ contains "almost" the restriction of $\mathbb{P}_{k}(P)$ to $\gamma$, i.e. that there exists a small function $\varepsilon$ which lives on $\gamma$ such that

$$
p_{k \mid \gamma} \notin V_{k}(\gamma) \quad \text { but } \quad p_{k \mid \gamma}+\varepsilon \in V_{k}(\gamma)
$$

Consider the virtual function $w_{h} \in V_{k}(P)$ that solves:

- $w_{h}=p_{k}+\varepsilon$ on $\gamma$
- $w_{h}=p_{k}$ on the other (straight) edges $e$;
- $\Delta w_{h}=\Delta p_{k} \in \mathbb{P}_{k-2}(P)$ in $P$.


## Does $V_{k}(P)$ still contain polynomials of degree $k$ ?

Then the difference $\left(w_{h}-p_{k}\right)$ solves the problem

- $\left(w_{h}-p_{k}\right)=\varepsilon$ on $\gamma$
- $\left(w_{h}-p_{k}\right)=0$ on the other (straight) edges;
- $\Delta\left(w_{h}-p_{k}\right)=0$ in $P$.

Hence

$$
w_{h}-p_{k}=O(\varepsilon) \text { on } \gamma \Longrightarrow w_{h}=p_{k}+O(\varepsilon) \text { in } P \quad \text { in the right norms. }
$$

In other words:

- if the space $V_{k}(\gamma)$ on the curved edge "approximates well" the restrictions of polynomials,
- then the virtual space $V_{k}(P)$ "approximates well" the polynomials "automatically".


## The "rough" and the "phony" way

So far we have explored two quite different ways for the definition of the space $V_{k}(\gamma)$ : the unofficial names are "rough" and "phony" ("rozzi" and "farlocchi" in italian).

- The "rough" way. In this case, we assume that the curve $\gamma$ is parametrized and we define $V_{k}(\gamma)$ as the functions which are polynomials of degree $k$ in the parameter.
- The "phony" way. Here we bite the bullet and define $V_{k}(\gamma)$ directly as the restrictions to $\gamma$ of all polynomials in two variables: $V_{k}(\gamma)=P_{k}(\gamma)$.

It is clear that if the edge $\gamma$ happens to be straight, in both cases we are back to the space of polynomials in one variable (unless you are crazy and do not parametrize the edge with an affine map!).

## The rough way



We then assume the the edge $\gamma$ is given in parametric form:

$$
t \mapsto \gamma(t)=(x(t), y(t)), \quad t \in[a, b] .
$$

## The rough way

For each parametrized curved edge $t \mapsto \gamma(t)$, we define the edge space $V_{k}(\gamma)$ as

$$
\begin{aligned}
& V_{k}(\gamma)=\{w: \gamma \rightarrow \mathbb{R}: \\
&\left.w(\gamma(t))=v_{h}(x(t), y(t)) \text { is a polynomial of degree } k\right\}
\end{aligned}
$$

It is clear that if $\gamma$ is straight, its parametrization is linear (affine) and we recover the standard VEM space.

## The degrees of freedom are the same as before

- pointwise values at the vertices;
- pointwise values at $k-1$ "equispaced" points on the edge (equispaced with respect to $t \in[a, b]$ for a curved edge);
- moments against scaled monomials up to degree $k-2$ inside.


Hence the dimension of $V_{k}(P)$ is the same as before.

## The $\Pi_{k}^{\nabla}$ projector onto $\mathbb{P}_{k}(P)$ can still be computed

We can still define the projector $\Pi_{k}^{\nabla}: V_{k}(P) \rightarrow \mathbb{P}_{k}(P)$.
As before, the projector $\Pi_{k}^{\nabla}$ is orthogonal with respect to the energy scalar product in $H^{1}(P)$ :

$$
\begin{gathered}
\int_{P} \nabla\left[\Pi_{k}^{\nabla} v_{h}\right] \cdot \nabla p_{k} \mathrm{~d} x=\int_{P} \nabla v_{h} \cdot \nabla p_{k} \mathrm{~d} x \quad \text { for all } p_{k} \in \mathbb{P}_{k}(P) \\
\int_{\partial P} \Pi_{k}^{\nabla} v_{h} \mathrm{~d} x=\int_{\partial P} v_{h} \mathrm{~d} x \quad \text { to fix the constant function }
\end{gathered}
$$

We show now that $\Pi_{k}^{\nabla} v_{h}$ is computable if we know the degrees of freedom of $v_{h}$.

## The $\Pi_{k}^{\nabla}$ projector onto $\mathbb{P}_{k}(P)$ can still be computed

We only need to check that the right hand side

$$
\int_{P} \nabla v_{h} \cdot \nabla p_{k} \mathrm{~d} x
$$

is computable from the degrees of freedom of $v_{h}$. Integrating by parts:

$$
\int_{P} \nabla v_{h} \cdot \nabla p_{k} \mathrm{~d} x=-\int_{P} v_{h} \Delta p_{k} \mathrm{~d} x+\int_{\partial P} v_{h}\left(\nabla p_{k} \cdot \boldsymbol{n}\right) \mathrm{d} s
$$

and the result follows by observing that

- $\Delta p_{k} \in \mathbb{P}_{k-2}(P)$ which implies that $\int_{P} v_{h} \Delta p_{k} \mathrm{~d} x$ is computable,
- $v_{h}$ is a known polynomial of degree $k$ on each straight edge $e$,
- $v_{h}$ is a known function in $V_{k}(\gamma)$ on each curved edge $\gamma$,
- $\nabla p_{k}$ is a known polynomial of degree $k-1$ on each edge.


## Issues

- The local VEM space $V_{k}(P)$ does not contain polynomials, since the corresponding edge space $V_{k}(\gamma)$ does not contain the restriction of polynomials in two variables. Hence the method will not pass the patch test.
- We need to integrate polynomials on curvilinear polygons and (known) functions on curves; of course in general we will not be able to compute the integrals exactly. Hence, the degree of precision of the quadrature formulas will come into play.


## Does it work??????

Roughly speaking, if a local space $V_{k}(P)$ contains all polynomials up to degree $k$, and we have a "regular" mesh sequence with $h$ going to zero, a Gakerkin method converges with the right order.

We can miss the polynomials, but not too much... otherwise convergence is lost.

The easiest case is when the curved edges come from the boundary of the domain or from an internal, fixed interface, which is the case covered by the isoparametric classical finite elements.

## Does it work??????



## Does it work??????



## Does it work??????




When the elements get smaller, they also get flatter. Hence for the rough VEM, the gap between polynomials and $V_{k}(P)$ diminishes as $h \longrightarrow 0$ and the method works.

## Main result

In the paper The Virtual Element Method with Curved Edges, ESAIM: M2AN 53 (2019), pp. 375-404 (also arXiv:1711.04306v2) together with L. Beirão da Veiga and G. Vacca we have proved that for the "fixed curved boundary" or the "fixed curved interface" problem, the Virtual Element Method described above converges with optimal rates, provided all integrals are computed without error.

This is the extension to VEM of the results that hold for the classical isoparametric Finite Elements.

From the numerical experiments it is apparent that it is enough to compute all integrals with the same precision necessary for the straight edge case. In other words, the bottom line is:
take into account the curved geometry but reuse the code for the straight edge case.

## General curved edges - not flattening

Jigsaw mesh


In this case we don't have convergence!

## Polynomial edges

Idea: we can raise the degree of the VEM on the edges independently of the accuracy $k$. The VEM local space is still well defined.

Assume that the curve describing the curved edge $\gamma$ is a polynomial of degree $d$ :

$$
\gamma(t)=(x(t), y(t)), \quad x, y \in \mathbb{P}_{d}(\mathbb{R})
$$



## The "enhanced rough" way

Now if $p_{k}(x, y)$ is a polynomial of degree $k$ in two variables, the function

$$
p_{k} \circ \gamma: t \mapsto p_{k}(x(t), y(t))
$$

is a polynomial of degree $k d$ in $t$.
Hence, in order to include all polynomials of degree $k$ in the local VEM space $V_{k}(P)$, we change the definition of $V_{k}(\gamma)$ as follows:

- for each curved edge $\gamma(t)$ which is a polynomial of degree $d$,

$$
v_{h} \in V_{k}(\gamma) \quad \text { means } \quad v_{h} \circ \gamma \in \mathbb{P}_{k d}(\mathbb{R})
$$

In other words, the function $t \mapsto v_{h}(x(t), y(t))$ must a polynomial of degree $k d$ in the paramatrization parameter $t$.

- for each straight edge $e, v_{h \mid e} \in \mathbb{P}_{k}(e)$ (as before)
- $\Delta v_{h} \in \mathbb{P}_{k-2}(P)$ (as before)


## The "enhanced rough" way

Jigsaw mesh, enhanced rough VEM


$$
\begin{gathered}
k=3, N d o f=11160, h_{\text {mean }}=6.736 \mathrm{e}-02 \\
L^{2} \text { error }=1.918 \mathrm{e}-05, H^{1} \text { error }=4.356 \mathrm{e}-04 \\
\text { edge error: } I^{2}=1.637 \mathrm{e}-05, \max =4.749 \mathrm{e}-05
\end{gathered}
$$


$\gamma=$ cubic polynomial, degree $3 \times 3=9$ on each edge

## The "enhanced rough" way

Jigsaw mesh, enhanced rough VEM


## The phony way

If we want to keep polynomials inside our local VEM spaces, we can simply define the edge space as the restriction to the edge of polynomials of degree $k$ in two variables, i.e.

$$
V_{k}(\gamma)=P_{k}(\gamma)
$$

If we proceed in this way, as we observed before we ensure that polynomials are included in $V_{k}(P)$, whatever the shape of $\gamma$.

This seems very natural, maybe a little bit expensive, but we don't need any parametrization - at least for the definition.

The real problem is the dimension of the space $P_{k}(\gamma)$ !

## Dimension of $P_{k}(\gamma)$

If $\gamma$ is a "generic" curve (i.e. not the zeros of some polynomial) then

$$
\operatorname{dim} \mathbb{P}_{k}(\gamma)=\operatorname{dim} \mathbb{P}_{k}\left(\mathbb{R}^{2}\right)=\frac{(k+1)(k+2)}{2}
$$

However, if $\gamma$ is straight then

$$
\operatorname{dim} \mathbb{P}_{k}(\gamma)=\operatorname{dim} \mathbb{P}_{k}\left(\mathbb{R}^{1}\right)=k+1
$$

and all intermediate case are possible!
The situation is even worse, since the dimension of $\mathbb{P}_{k}(\gamma)$ is unstable with respect to small changes in $\gamma$, and I think it would be really hard to decide in the code case by case the "numerical" dimension of $\mathbb{P}_{k}(\gamma)$.

After long discussions, our point of view is the following: we don't want to check the dimension of $\mathbb{P}_{k}(\gamma)$. We want a method that works in ALL cases, even when $\gamma$ is a segment. In other words: we don't want to care.

## Degrees of freedom for $P_{k}(\gamma)$

Since $P_{k}(\gamma)$ are the restrictions to $\gamma$ of polynomials in two variables up to degree $k$, we describe $P_{k}(\gamma)$ with the polynomials in two variables up to degree $k$. Hence we start from a mesh with a curved edge $\gamma$ :


## Degrees of freedom for $P_{k}(\gamma)$

For the curved edge $\gamma$ we construct the Lagrange nodes for polynomials in two variables up to degree $k$ :


We identify a function in $V_{k}(\gamma)$ by its pointwise values in the Lagrange nodes: the "phony" dofs.

## Degrees of freedom for $P_{k}(\gamma)$

The phony dofs at the vertices are special because they are also "true" dofs, connecting the two worlds:


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## Degrees of freedom for $V_{k}(P)$

We go back now to our polygon $P$ and the local VEM space $V_{k}(P)$ :


There are 8 phony dofs $\{*\}$ and 22 true dofs $\{\bullet, \square\}$. If the edge $\gamma$ becomes straight, then 6 of the phony dofs become useless. All intermediate cases are possible.

## Degrees of freedom for $V_{k}(P)$

Hence $\operatorname{dim} V_{k}(P)$ is between 24 and 30 depending on $\gamma$.
Since we don't want to decide what kind of curve $\gamma$ is, we keep 30 parameters to describe $V_{k}(P)$.

Note that we do not assume that they are degrees of freedom.
Given an array $\boldsymbol{V}_{\boldsymbol{h}}$ in $\mathbb{R}^{30}$, we define the associated function $v_{h} \in V_{k}(P)$ in the following way:

- Take the 10 values of $\boldsymbol{V}_{\boldsymbol{h}}$ corresponding to the 8 phony dofs and the 2 true dofs on the vertices of $\gamma$;
- build the corresponding interpolation polynomial $p_{k}$ in two variables (recall that in the example $k=3$ );
- define $v_{h}$ on $\gamma$ as the restriction of $p_{k}$ on $\gamma$.
- Use the true edge dofs of $\boldsymbol{V}_{\boldsymbol{h}}$ to determine $v_{h}$ on the other edges.
- Use the true internal dofs of $\boldsymbol{V}_{\boldsymbol{h}}$ to determine $v_{h}$ inside.


## Degrees of freedom for $V_{k}(P)$

The reverse path would be:

- Take a function $v_{h} \in V_{k}(P)$;
- restrict $v_{h}$ on the edge $\gamma$;
- from the restriction of $v_{h}$ to $\gamma$ identify a polynomial $p_{k}$ in two variables;
- the array $\boldsymbol{V}_{\boldsymbol{h}} \in \mathbb{R}^{30}$ is given by
- the true dofs of $v_{h}$ for straight edges and in the interior of $P$;
- the value of $p_{k}$ at the phony dofs

As we have already seen, this path is unstable.

$$
\begin{aligned}
& \boldsymbol{V}_{\boldsymbol{h}} \Longrightarrow v_{h} \text { is OK; } \\
& v_{h} \Longrightarrow \boldsymbol{V}_{\boldsymbol{h}} \text { is forbidden. }
\end{aligned}
$$

## Degrees of freedom for $V_{k}(P)$

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- the array $\boldsymbol{V}_{\boldsymbol{h}} \in \mathbb{R}^{30}$ is given by
- the true dofs of $v_{h}$ for straight edges and in the interior of $P$;
- the value of $p_{k}$ at the phony dofs

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\end{aligned}
$$

## Degrees of freedom for $\mathbb{P}_{k}(P)$

We have seen that we can go from the parameters in $\mathbb{R}^{30}$ to a function but not viceversa.

There is a notable exception: we can reverse the path when the function is
a polynomial $p_{k}$.
In fact, regardless of the shape of $\gamma$, we can always define an array $\boldsymbol{P}_{\boldsymbol{k}}$ associated with a polynomial $p_{k}$ by simply computing the value of $p_{k}$ at the phony dofs, simply observing that polynomials are globally defined.
$\boldsymbol{P}_{\boldsymbol{k}} \Longleftrightarrow p_{k}$ is OK for polynomials

$$
\boldsymbol{P}_{\boldsymbol{k}} \Longrightarrow p_{k} \Longrightarrow \boldsymbol{P}_{\boldsymbol{k}}
$$

## The phony method

The simplest VEM for the Laplace operator defines the local approximate bilinear form on $P$ in the following way:

$$
\begin{aligned}
& a_{h}^{P}\left(u_{h}, v_{h}\right)= \\
& \qquad \int_{P} \nabla \Pi_{k}^{\nabla} u_{h} \cdot \nabla \Pi_{k}^{\nabla} v_{h} \mathrm{~d} x+\sum_{\ell} \operatorname{dof}_{\ell}\left(u_{h}-\Pi_{k}^{\nabla} u_{h}\right) \operatorname{dof}_{\ell}\left(v_{h}-\Pi_{k}^{\nabla} v_{h}\right)
\end{aligned}
$$

where the index $\ell$ in the summation runs across all dofs.
This is the so-called "dofi-dofi" stabilization.

## The phony method

Since our unknown will be the array $\boldsymbol{U}_{\boldsymbol{h}}$, we need to find a related approximate bilinear form defined on arrays:

$$
\begin{aligned}
& a_{h}^{P}\left(\boldsymbol{U}_{\boldsymbol{h}}, \boldsymbol{V}_{\boldsymbol{h}}\right)= \\
& \qquad \int_{P} \nabla \Pi_{k}^{\nabla} \boldsymbol{U}_{\boldsymbol{h}} \cdot \nabla \Pi_{k}^{\nabla} \boldsymbol{V}_{\boldsymbol{h}} \mathrm{d} x+\sum_{\ell} \operatorname{dof}_{\ell}\left(\boldsymbol{U}_{\boldsymbol{h}}-\Pi_{k}^{\nabla} \boldsymbol{U}_{\boldsymbol{h}}\right) \operatorname{dof}_{\ell}\left(\boldsymbol{V}_{\boldsymbol{h}}-\Pi_{k}^{\nabla} \boldsymbol{V}_{\boldsymbol{h}}\right)
\end{aligned}
$$

Now:

- $\Pi_{k}^{\nabla} \boldsymbol{V}_{\boldsymbol{h}}$ is defined through the projection of the associated function $v_{h}$ :

$$
\boldsymbol{V}_{\boldsymbol{h}} \Longrightarrow v_{h} \Longrightarrow \Pi_{k}^{\nabla} v_{h}
$$

- $\operatorname{dof}_{\ell}\left(\boldsymbol{V}_{\boldsymbol{h}}\right)$ is simply the $\ell$-th component of the array $\boldsymbol{V}_{\boldsymbol{h}}$;
- $\Pi_{k}^{\nabla} \boldsymbol{V}_{\boldsymbol{h}}=\Pi_{k}^{\nabla} v_{h}$ is a polynomial of degree $k$, and we have seen that we can associate to it an array of dofs in a consistent way.


## Local stiffness matrix in the singular case

When the curve $\gamma$ "degenerates", the consistency part becomes singular but the stability takes care of it:


## Local stiffness matrix in the singular case

We see now what happens to the local stiffness matrix in this case.

- Suppose that $\gamma$ is straight and take the array $\boldsymbol{\Phi}_{\boldsymbol{i}}$ which is 1 only on the useless phony dof $i$ and 0 on the other dofs (true and phony).
- The corresponding basis function $\varphi_{i}$ is identically zero, and the same obviously happens for $\Pi_{k}^{\nabla} \varphi_{i}$.
- Hence the consistency term is all zero:

$$
\int_{P} \nabla \Pi_{k}^{\nabla} \boldsymbol{\Phi}_{\boldsymbol{j}} \cdot \nabla \Pi_{k}^{\nabla} \boldsymbol{\Phi}_{\boldsymbol{i}} \mathrm{d} x=0 \text { for all } j
$$

- The stability term reduces to

$$
\begin{aligned}
\sum_{\ell} \operatorname{dof}_{\ell}\left(\boldsymbol{\Phi}_{\boldsymbol{j}}-\Pi_{k}^{\nabla} \boldsymbol{\Phi}_{\boldsymbol{j}}\right) \operatorname{dof}_{\ell}\left(\boldsymbol{\Phi}_{\boldsymbol{i}}-\Pi_{k}^{\nabla} \boldsymbol{\Phi}_{\boldsymbol{i}}\right)= \\
\sum_{\ell} \operatorname{dof}_{\ell}\left(\boldsymbol{\Phi}_{\boldsymbol{j}}-\Pi_{k}^{\nabla} \boldsymbol{\Phi}_{\boldsymbol{j}}\right) \delta_{\ell i}=\delta_{j i}-\operatorname{dof}_{i}\left(\Pi_{k}^{\nabla} \boldsymbol{\Phi}_{\boldsymbol{j}}\right)
\end{aligned}
$$

and if $j$ is also a useless phony dof, then we get the identity matrix.

## Local stiffness matrix in the singular case

Local VEM stiffness matrix


## Numerical experiments



The curved edges degenerate to straight edges.

## Numerical experiments - patch test $k=3$

$$
\begin{gathered}
k=3, \text { Ndof }=488, h_{\text {mean }}=3.047 \mathrm{e}-01 \\
L^{2} \text { error }=5.015 \mathrm{e}-14, \mathrm{H}^{1} \text { error }=7.474 \mathrm{e}-14 \\
\text { edge error: } \mathrm{I}^{2}=4.233 \mathrm{e}-14, \max =7.572 \mathrm{e}-16
\end{gathered}
$$

$$
\begin{gathered}
k=3, \text { Ndof }=488, h_{\text {mean }}=3.047 \mathrm{e}-01 \\
\mathrm{~L}^{2} \text { error }=2.825 \mathrm{e}-14, \mathrm{H}^{1} \text { error }=4.673 \mathrm{e}-14 \\
\text { edge error: } \mathrm{I}^{2}=2.466 \mathrm{e}-14, \max =7.572 \mathrm{e}-16
\end{gathered}
$$



$\gamma$ straight

The solution doesn't change when the curved edges degenerate to straight edges.

## Numerical experiments - Poisson equation $k=3$

$$
\begin{gathered}
k=3, N d o f=2670, h_{\text {mean }}=1.322 \mathrm{e}-01 \\
L^{2} \text { error }=7.099 e-04, H^{1} \text { error }=5.058 \mathrm{e}-03 \\
\text { edge error: } I^{2}=5.298 \mathrm{e}-04, \max =6.556 \mathrm{e}-05
\end{gathered}
$$


$\gamma=\sin$

$$
\begin{gathered}
\mathrm{k}=3, \mathrm{Ndof}=2670, \mathrm{~h}_{\text {mean }}=1.322 \mathrm{e}-01 \\
\mathrm{~L}^{2} \text { error }=7.047 \mathrm{e}-04, \mathrm{H}^{1} \text { error }=5.013 \mathrm{e}-03 \\
\text { edge error: } \mathrm{I}^{2}=5.125 \mathrm{e}-04, \max =6.556 \mathrm{e}-05
\end{gathered}
$$

$\gamma$ straight
The solution doesn't change when the curved edges degenerate to straight edges.

## Numerical experiments - convergence



$$
\gamma=\text { quadratic }
$$

Here the curved edges $\gamma$ are quadratic polynomials and $k=3$. Hence the dimension of $V_{k}(\gamma)$ is not equal to the dimension of polynomials in two variables, and there are phony dofs that are useless. The methods works in any case, without need to check!

## Numerical experiments - convergence $k=3$



$\gamma=$ quadratic
In all the experiments I have overkilled the integration over polygons with curved edges. The minimum degree of precision of the quadrature formulas to be used in order to preserve optimal convergence is unknown.

## Conclusion and perspectives

- There is a lot of work to do!
- The classical VEM machinery continue to work (enhancing for $L^{2}$ projectors, variable coefficients,...).
- The methods presented could be of interest also for triangular and quadrilateral FEM meshes with curved edges.
- Integration over curved polygons and polyhedra must be investigated.
- Extension to three dimensions?


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## Thanks for your attention!

