

# Dualité de Galois-Grothendieck et dualité tannakienne

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- 1 Galois-Grothendieck duality and Tannaka duality
- 2 Hopf adjunctions and Hopf comonads
- 3 Main result
- 4 Applications

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This follows from

### Theorem

Let  $C$  be a category and  $\omega : C \rightarrow \text{set}$  be a functor, and assume :

- ①  $C$  has finite limits and colimits, and  $\omega$  is exact ;
- ② in  $C$  any morphism factorizes as  $\text{mono} \circ \text{epi}$ , and epis are strict ;
- ③  $\omega$  is conservative ;
- ④ in  $C$  monos are summands.

Then  $C \cong G - \text{set}$ , where  $G$  is the profinite group  $\text{Aut}(\omega)$ .

Tannaka theory in its algebraic form (Saavedra Rivano's thesis in 1972, Deligne-Milne, Deligne...).

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**Neutral case.** *Let  $C$  be a symmetric tensor category over a field  $\mathbb{k}$ ,  $\omega : C \rightarrow \text{Vect}_{\mathbb{k}}$  a symmetric fibre functor, that is, a strong monoidal symmetric  $\mathbb{k}$ -linear exact functor.*



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2) **General case.** *Still works if you replace  $\text{Vect}_{\mathbb{k}}$  with  $\text{Mod}_B$ , with  $B$  commutative  $\mathbb{k}$ -algebra  $\neq 0$ . Then  $G$  is an affine groupoid with base  $\text{Spec}B$ .*

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- Non-commutative case, where  $\mathcal{C}$  is no longer symmetric  $\rightsquigarrow G$  is a 'quantum' Hopf algebroid [Maltsiniotis, B]
- Fully non-commutative case,  $B$  non-commutative,  $\omega : \mathcal{C} \rightarrow_B \text{Bimod}_B \rightsquigarrow G$  is a Hopf algebroid in the sense of Takeuchi/Schauenburg

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Our first step to address this question will be to study (co)monoidal adjunction and introduce Hopf (co)monads which serve exactly that kind of purpose.

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A **monoidal adjunction** is an adjunction  $\omega \begin{array}{c} \mathcal{C} \\ \curvearrowright \\ \mathcal{B} \end{array} \begin{array}{l} \\ R \end{array}$  where  $\mathcal{C}, \mathcal{B}$  are monoidal and  $\omega$  is strong monoidal.

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It is **Hopf** if for all  $b \in \mathcal{B}, c \in \mathcal{C}$  the **fusion morphisms** are isomorphisms :

$$H_{b,c}^l : c \otimes Rb \rightarrow R(\omega c \otimes b) \quad \text{and} \quad H_{b,c}^r : Rb \otimes c \rightarrow R(b \otimes \omega c)$$

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### Example

$B$  bialgebra over  $\mathbb{k} \rightsquigarrow$  a monoidal adjunction  $U \begin{array}{c} \text{Comod } B \\ \leftarrow \quad \rightarrow \\ \text{Vect}_{\mathbb{k}} \end{array} \text{cofree}$



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It is Hopf iff  $B$  is a Hopf algebra (w. bijective antipode).

# Hopf comonads

A **monoidal comonad** [I. Moerdijk] on a monoidal category  $\mathcal{B}$  is a comonad  $(T, \Delta, \varepsilon)$  such that  $T$ ,  $\Delta$  and  $\varepsilon$  are monoidal, that is :

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$$\begin{array}{ll}
 T: \mathcal{B} \rightarrow \mathcal{B}, & \Delta: T \rightarrow T^2 \text{ (coproduct),} & \varepsilon: T \rightarrow 1_{\mathcal{B}} \text{ (counit)} \\
 \mu_{a,b}: Ta \otimes Tb \rightarrow T(a \otimes b), & & \eta: 1 \rightarrow T1
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The  $\mu - \Delta$  compatibility doesn't require a braiding :

$$\begin{array}{ccc}
 Ta \otimes Tb & \xrightarrow{\Delta_a \otimes \Delta_b} & T^2 a \otimes T^2 b \xrightarrow{\Delta_{Ta, Tb}} T(Ta \otimes Tb) \\
 \mu_{a,b} \downarrow & & \downarrow T\mu_{a,b} \\
 T(a \otimes b) & \xrightarrow{\Delta_{a \otimes b}} & T^2(a \otimes b)
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Define the (left and right) *fusion morphisms*

- $H_{a,b}^l = \mu_{a, Tb}(Ta \otimes \Delta_b): Ta \otimes Tb \rightarrow T(a \otimes Tb),$
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$T$  is a *Hopf comonad* if  $H^l$  and  $H^r$  are isomorphisms.

## Hopf comonads and Hopf monoidal adjunctions

$T$  comonad on  $\mathcal{B} \rightsquigarrow$  an adjunction  $U_T \left( \begin{array}{c} \mathcal{B}_T \\ \uparrow \\ \mathcal{B} \\ \downarrow \\ \mathcal{B} \end{array} \right) R_T$

where  $\mathcal{B}_T = \{T\text{-comodules } (b, \delta : b \rightarrow Tb)\}$ ,  $U_T(b, \delta) = b$  and  $R_T(b) = (Tb, \Delta_b)$ .

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## Hopf comonads and Hopf monoidal adjunctions

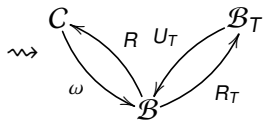
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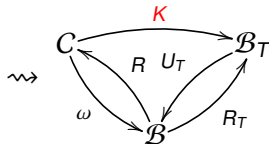
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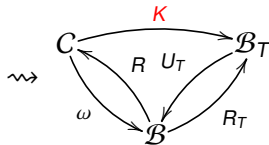
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the *comparison functor*  $K : c \mapsto (\omega c, \omega \eta_c)$  is strong monoidal

The adjunction  $(\omega, R)$  is *comonadic* if  $K$  is an equivalence.

# Outlook of General Theory of Hopf (co)monads

10/23

The notion of a Hopf comonad is not self-dual, unlike that of a Hopf algebra. The dual notion is that of Hopf monads. Many classical results of the theory of Hopf algebras extend to Hopf (co)monads [BV 2007, BVL 2011]

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- Applications to construction and comparison of quantum invariants (non-braided setting)

# Hopf comonads from Hopf algebras

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**Consequence for Tannaka duality** : the neutral case means that  $\omega$  has a monoidal section  $\rightsquigarrow$  the reconstructed Hopf structure is a Hopf algebra in  $\mathcal{Z}(\mathcal{B})$ .

## Hopf comonads as 'quantum groupoids'

Let  $R$  be a unitary ring  $\rightsquigarrow$  a monoidal category  $({}_R\text{Mod}_R, \otimes_{R,R}, R_R)$ .

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Hopf algebroids are non-commutative avatars of groupoids. Complicated axioms  $\rightsquigarrow$  a Hopf adjunction  $\rightsquigarrow$  a Hopf comonad (much easier to manipulate).

- 1 Galois-Grothendieck duality and Tannaka duality
- 2 Hopf adjunctions and Hopf comonads
- 3 Main result**
- 4 Applications

# Reconstruction without an adjoint

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The Hopf comonad  $T_{\omega}$  will play the rôle of  $\text{Aut}(\omega)$ .

# Comonadicity criterion

Let  $\omega : \mathcal{C} \rightarrow \mathcal{B}$  be a strong monoidal functor, with  $\mathcal{B}$  having small filtered colimits which are exact and preserved by  $\otimes_{\mathcal{B}}$ . Assume that

- 1  $\mathcal{C}$  has finite limits and colimits and  $\omega$  is exact ;
- 2  $\mathcal{C}$  has mono-epi factorizations and has strict monos ;
- 3  $\mathcal{C}$  is coartinian ;
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Then

- a)  $\omega$  extends uniquely to a strong monoidal functor  $\Omega : \text{Ind}\mathcal{C} \rightarrow \mathcal{B}$  which preserves filtered colimits and has a right adjoint  $R$  ;
- b)  $\Omega$  is comonadic with comonad  $T = T_{\omega} = \Omega R$ , so that  $\text{Ind}\mathcal{C} \cong_{\otimes} \mathcal{B}_T$  ;
- c) Moreover, if  $\mathcal{B}^f \subset \mathcal{B}$  is a full monoidal subcategory containing  $\omega(\mathcal{C})$  and whose objects have finite type in  $\mathcal{B}$  (e. g.  $\mathcal{B} = \text{Ind}\mathcal{B}^f$ ), we have  $\mathcal{C} \cong_{\otimes} \mathcal{B}_T^f$ .

$$\begin{array}{ccc} \text{Ind}C & \xrightarrow{\simeq_{\otimes}} & \mathcal{B}_T \\ & \searrow \Omega & \swarrow \\ & \mathcal{B} & \end{array}$$

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If in addition

- 1  $C$  has duals or  $C$  and  $\omega$  are coclosed,
  - 2 for any mono  $i$  of  $C$ ,  $\omega(i)$  is a tensor-universal mono of  $\mathcal{B}$ ,
- then  $T$  is a Hopf comonad.

Note that, if  $\omega$  has a monoidal section, then the Hopf comonad  $T$  is co-augmented, so there exists a Hopf algebra  $(H, \sigma)$  in  $\mathcal{Z}(\mathcal{B})$  such that  $T = H \otimes_{\sigma}$  and  $C$  is the category of  $H$ -comodules in  $\mathcal{B}$ .

- 1 Galois-Grothendieck duality and Tannaka duality
- 2 Hopf adjunctions and Hopf comonads
- 3 Main result
- 4 Applications**
  - Recovering classical results
  - The hidden commutative algebra
  - Application to tensor functors



In its direct form the main result yields tannaka duality in the neutral case and in the general case (even non-commutative), with  $\mathcal{B} = \text{Vect}_{\mathbb{k}}$ ,  
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The first part of the theorem reconstructs a comonoidal monad  $T$  on  $\text{Proset}$ , which is in fact a bialgebra  $G$  because  $\text{set}$  is generated by the unit object and  $T$  preserves finite sums and filtering limits. Since we are in a cartesian category, bialgebras are just monoids.

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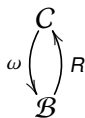
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One can refine the criterion for categories with enough duals.

# Hopf adjunctions revisited

19/23

Consider a **Hopf monoidal adjunction**



Theorem (BVL)

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## Theorem (BVL)

$A = R(\mathbb{1})$  is an algebra in  $\mathcal{C}$ , and comes with a canonical half-braiding  $\sigma$  which makes it a **commutative algebra in  $\mathcal{Z}(\mathcal{C})$**  called the *induced central algebra*.

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The central commutative algebra generalizes Deligne’s trivializing algebra.

Let  $F : \mathcal{C} \rightarrow \mathcal{D}$  be a tensor functor between tensor categories. Then  $\text{Ind}F : \text{Ind}\mathcal{C} \rightarrow \text{Ind}\mathcal{B}$  has a right adjoint  $R$ .

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## Theorem

If  $R$  is faithful exact ( $F$  dominant), with induced central algebra  $(A, \sigma)$ , we have a commutative diagram

$$\begin{array}{ccc}
 C & \xrightarrow{F_A} & C_{A,\sigma} \\
 & \searrow F & \nearrow \simeq_{\otimes} \\
 & \mathcal{D} &
 \end{array}$$

where  $C_A$  is the category of  $A$ -modules 'of finite type' and  $F_A$  is the tensor functor  $X \mapsto A \otimes X$ .

If  $\mathcal{D} = \text{vect}_{\mathbb{k}}$  and  $C, F$  are symmetric, then  $A$  is Deligne's trivializing algebra.

# Existence of fibre functors

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## Theorem [BLV]

Let  $\mathcal{C}$  be a finite tensor category. Then there exists a finite dimensional  $\mathbb{k}$  algebra  $B$  and a fibre functor  $\omega : \mathcal{C} \rightarrow_B \text{Bimod}_B$ . Hence  $\mathcal{C} \cong_{\otimes} \text{rep}A$  for a certain Hopf algebroid with base  $B$ .



## 24. More on Hopf monads

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**BV1.** *Hopf Diagrams and Quantum Invariants*, AGT **5** (2005) 1677-1710.

Where Hopf diagram are introduced as a means for computing the Reshetikhin-Turaev invariant in terms of the coend of a ribbon category and its structural morphisms.

**BV2.** *Hopf Monads*, Advances in Math. **215** (2007), 679-733.

Where the notion of Hopf monad is introduced, and several fundamental results of the theory of finite dimensional Hopf algebras are extended thereto.

**BV3.** *Categorical Centers and Reshetikhin-Turaev Invariants*, Acta Mathematica Vietnamica **33** 3, 255-279

Where the coend of the center of a fusion spherical category over a ring is described, the modularity of the center, proven, and the corresponding Reshetikhin-Turaev invariant, constructed.

**BV4.** *Quantum Double of Hopf monads and Categorical Centers*, arXiv :0812.2443, to appear in Transactions of the American Mathematical Society (2010)

Where the general theory of centralizers and doubles of Hopf monads is expounded.

**BLV.** *Hopf Monads on Monoidal Categories*, arXiv :1003.1920.

Where Hopf monads are defined anew in the monoidal world

**BN.** *Exact sequences of tensor categories*, arXiv :1006.0569.

See also : <http://www.math.univ-montp2.fr/~bruguieres/recherche.html>