

On Kazhdan-Lusztig cells in type B

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Weyl group

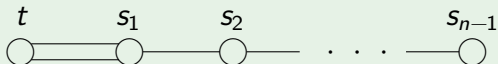
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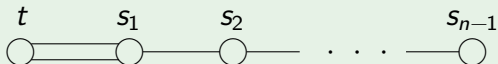
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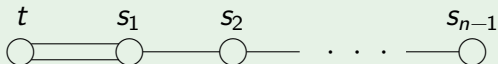


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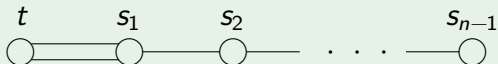


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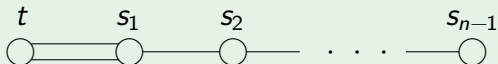


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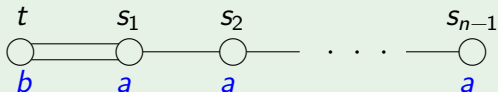
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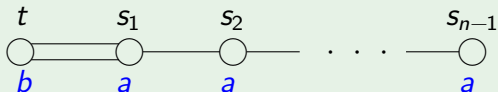
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- For simplification, $a, b > 0$.

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- **Involution:** $\overline{e^\gamma} = e^{-\gamma}$, $\overline{T_w} = T_{w^{-1}}^{-1}$

Theorem (Kazhdan-Lusztig 1979, Lusztig 1983)

If $w \in W_n$, there exists a unique $C_w \in \mathcal{H}_n$ such that

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$$C_{s_1 t s_1} = T_{s_1 t s_1} + q^{-1}(T_{s_1 t} + T_{t s_1}) + q^{-2} T_t$$
$$+ Q^{-1} q^{-1} (T_{s_1} + q^{-1}) \times \begin{cases} (1 + q^2) & \text{if } b > a, \\ 1 & \text{if } b = a, \\ (1 - Q^2) & \text{if } b < a \end{cases}$$

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- By construction, $I_{\leq_L \mathcal{C}}$ and $I_{<_L \mathcal{C}}$ are left ideals of \mathcal{H}_n and $V_{\mathcal{C}}$ is a left \mathcal{H}_n -module: $V_{\mathcal{C}}$ is called the **left cell representation** associated to \mathcal{C} .

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$$x \leq_L y \iff x^{-1} \leq_R y^{-1}$$

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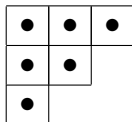
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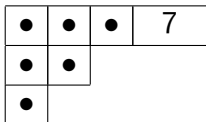
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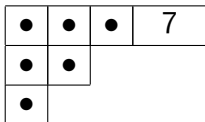
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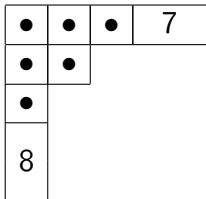
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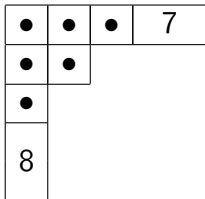
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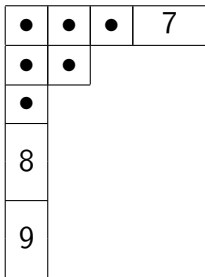
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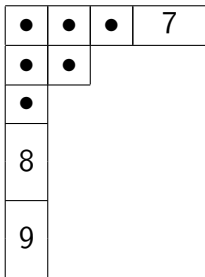
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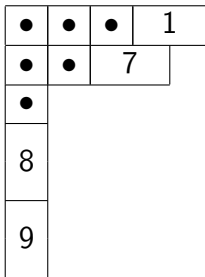
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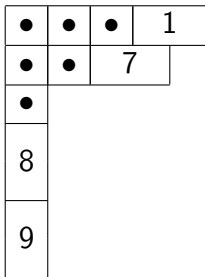
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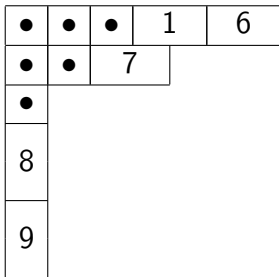
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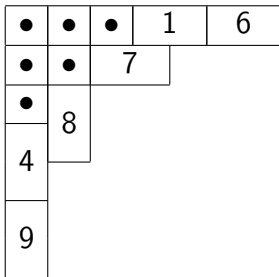
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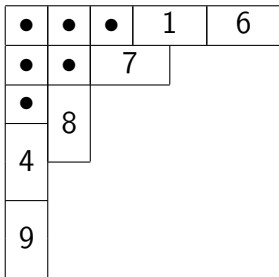
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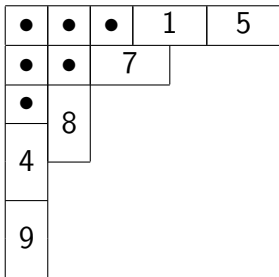
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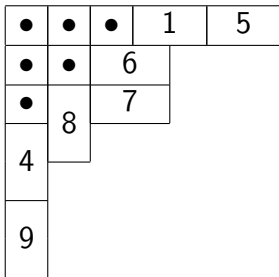
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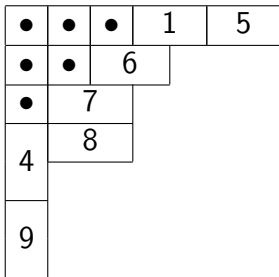
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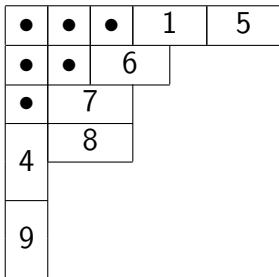
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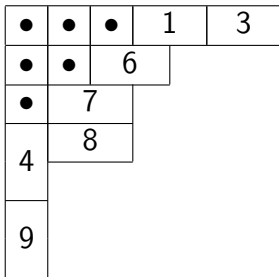
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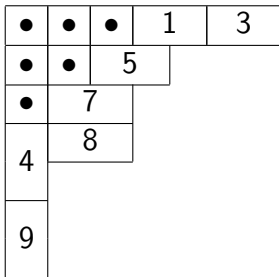
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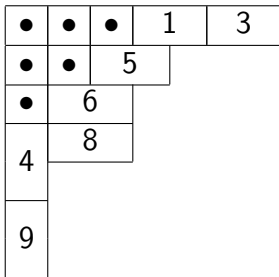
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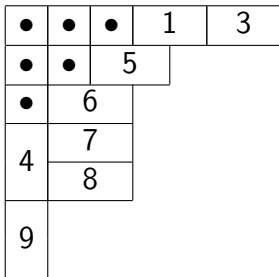
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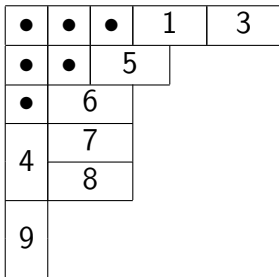
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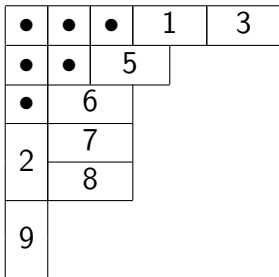
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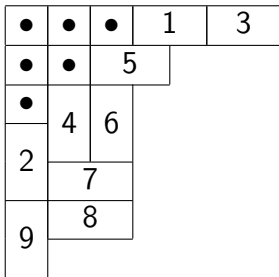
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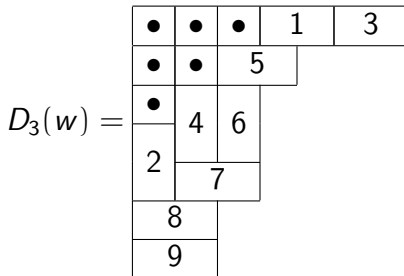
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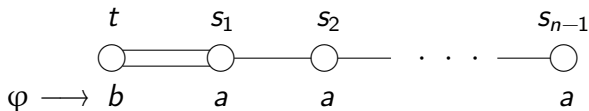
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Theorem (Garfinkle, van Leeuwen)

$$\begin{aligned} W_n &\xrightarrow{\sim} SDT_r^{(2)}(n) \\ w &\mapsto (D_r(w), D_r(w^{-1})) \end{aligned}$$

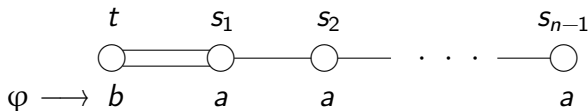
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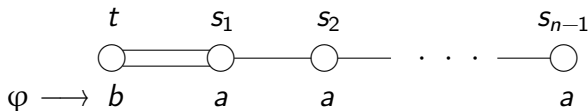
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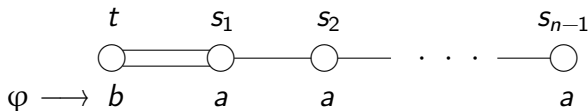
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REMARK - It should be possible to remove the hypothesis on b in the previous corollary using work of Pietraho.

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Lemma (1)

Let $w \in W_n$, let $i \in I_{n-1}^+$ and assume that one of the following holds:

- (1) $i \geq 2$ and $w(i) < w(i-1) < w(i+1)$,
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Then $w \sim_R ws_i$.

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Lemma 2 is implied by

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Let $l \in \{1, \dots, n-1\}$ and assume that $b \geq (n-1)a$. Then the coefficient of $C_{a_l \sigma_{[l+1, n]}}$ in $C_t C_{s_{n-1} \cdots s_{l+1} s_l s_1 s_2 \cdots s_{l-1} a_l \sigma_{[l+1, n]}}$ is non-zero!

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