

Multivariate growth and cogrowth

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Abstract

We investigate a multivariate growth series $\Gamma_L(\mathbf{z})$, $\mathbf{z} \in \mathbb{C}^d$ associated with a regular language L over an alphabet of cardinality d . Our focus is on languages coming from subgroups of the free group and from subshifts of finite type. We develop a mechanism for computing the rate of growth $\varphi_L(\mathbf{r})$ of L in the direction $\mathbf{r} \in \mathbb{R}^d$. Using the concave growth condition (CG) introduced by the second author in [18] and the results of Convex Analysis we represent $\psi_L(\mathbf{r}) = \log(\varphi_L(\mathbf{r}))$ as a support function of a convex set that is a Relog image of the domain of absolute convergence of $\Gamma_L(\mathbf{z})$. This allows us to compute $\psi_L(\mathbf{r})$ in some important cases, like a Fibonacci language or a language of freely reduced words representing elements of a free group F_2 . Also we show that the methods of the Large deviation theory can be used as an alternative approach. Finally, we suggest some open problems directed on the possibility of extensions of the results of the first author from [13] on multivariate cogrowth.

Keywords: Growth, Cogrowth, Regular language, Multivariate growth exponent, Free group, Fibonacci subshift, Subshift of finite type, Large deviations principle

Mathematics Subject Classification – MSC2020: 20E05, 20F69, 05A05, 05A15, 05A16, 60F10, 68Q45

1 Introduction

The study of asymptotic properties of a sequence $\{\gamma_n\}_{n \geq 0}$ of real (or complex) numbers is related to study of analytic properties of the function $\Gamma(z)$ presented by a power series $\sum_{n=0}^{\infty} \gamma_n z^n$. This includes inspection of singularities on the border

of the domain of (absolute) convergence of $\Gamma(z)$ and local behavior of $\Gamma(z)$ in their neighborhood. If $\Gamma(z)$ is a rational function i.e.

$$\Gamma(z) = \frac{G(z)}{H(z)}, \quad G(z), H(z) \in \mathbb{C}[z]$$

then all needed information about the coefficients γ_n can be gained from the polynomials $G(z)$ and $H(z)$.

In algebra, and especially in modern group theory, there are many notions and concepts that attached to the algebraic object a sequence $\{\gamma_n\}_{n \geq 0}$. This include growth, cogrowth, subgroup growth etc. Recall that given a finitely generated group G with a system of generators S , one can consider a function

$$\gamma_n = \#\{g \in G : |g| = n\},$$

where $n \in \mathbb{N}$ and $|g|$ is the length of the element g with respect to S . If the pair (G, S) has a regular geodesic normal form (in other terminology a rational cross section [7]) then the power series

$$\Gamma(z) = \sum_{n=0}^{\infty} \gamma_n z^n$$

represents a rational function and the asymptotic of γ_n is either polynomial or exponential. There are many groups (for instance groups of intermediate growth constructed in [8]) for which $\Gamma(z)$ is irrational for any system of generators and the study of asymptotic properties of $\{\gamma_n\}_n^\infty$ becomes much more complicated.

Now let F_m be a free group of rank m . Every m generated group G can be presented as a quotient F_m/N , for a suitable normal subgroup $N \triangleleft F_m$. Let $A = \{a_1, \dots, a_m\}$ be a basis of F_m . Elements of F_m are presented by freely reduced words over alphabet $\Sigma = \{a_1, \dots, a_m, a_1^{-1}, \dots, a_m^{-1}\}$ and there is $2m(2m-1)^{n-1}$ such words of length $n \geq 1$. The function

$$\Gamma_{F_m}(z) = \sum_{n=0}^{\infty} 2m(2m-1)^{n-1} z^n = \frac{1+z}{1-(2m-1)z}$$

is a spherical growth function of F_m with respect to the basis S . Now let $H < F_m$ be a subgroup and H_n be the set of elements in H of length n with respect to generators $\{a_1, \dots, a_m\}$ of F_m . The sequence $\{|H_n|\}_n^\infty$ of cardinalities of these sets is a cogrowth sequence,

$$H(z) = \sum_{n=0}^{\infty} |H_n| z^n$$

is a cogrowth series, and

$$\alpha_H = \limsup_{n \rightarrow \infty} |H_n|^{\frac{1}{n}}$$

is a cogrowth. The range for α_H is $[1, 2m - 1]$ and the range of α_H when H is nontrivial and normal subgroup is $(\sqrt{2m - 1}, 2m - 1]$. (See [10, 11, 12, 13]). The spectral radius χ of the simple random walk on $G = F_m/N$, when $N \triangleleft F_m$ is related with α_N as

$$\chi = \frac{\sqrt{2m - 1}}{2m} \left(\frac{\sqrt{2m - 1}}{\alpha_N} + \frac{\alpha_N}{\sqrt{2m - 1}} \right)$$

and the group G is amenable if and only if $\alpha_N = 2m - 1$ (that is α_N takes its maximum possible value).

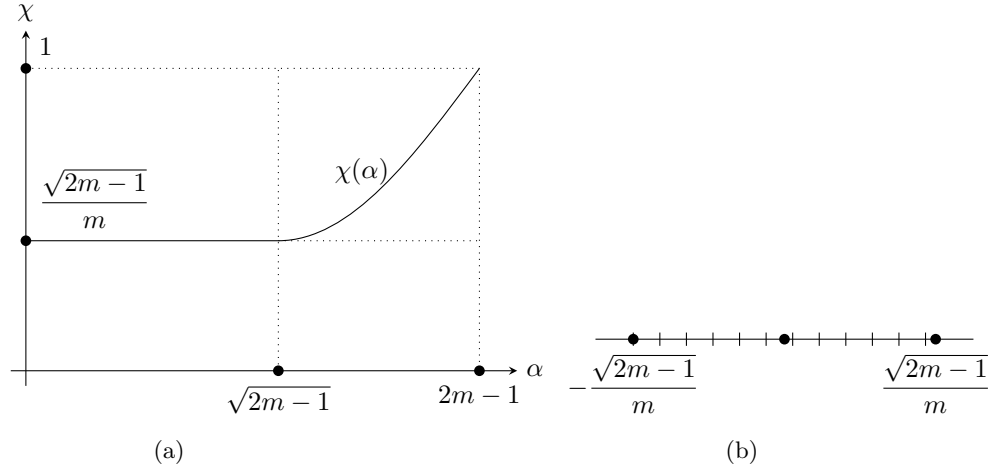


Figure 1: The graph of $\chi = \chi(\alpha)$ and the interval $\left[-\frac{\sqrt{2m - 1}}{m}, \frac{\sqrt{2m - 1}}{m} \right]$

In the case when $H < F_m$ is not a normal subgroup, one can consider a Schreier graph $\Lambda = \Lambda(F_m, H, \Sigma)$. Then the dependence on α_H of the spectral radius χ of a simple random walk on Γ is given by

$$\chi = \begin{cases} \frac{\sqrt{2m - 1}}{2m} \left(\frac{\sqrt{2m - 1}}{\alpha_H} + \frac{\alpha_H}{\sqrt{2m - 1}} \right) & \text{if } \alpha_H > \sqrt{2m - 1} \\ \frac{\sqrt{2m - 1}}{2m} & \text{if } \alpha_H \leq \sqrt{2m - 1} \end{cases}$$

(See Figure 1a). Hence, again the graph Γ is amenable if and only if $\alpha_H = 2m - 1$, while in the case when Γ is infinite, it is a Ramanujan graph if and only if $\alpha_H \leq \sqrt{2m - 1}$. The value $\frac{\sqrt{2m - 1}}{m}$ is a spectral radius of random walk on F_m computed by H. Kesten [14] and if the spectrum of the Laplacian operator of graph Γ minus two points set $\{-1, 1\}$ is a subset of the interval given by Figure 1b, then the graph is called Ramanujan. Observe that the analogue of the formula for χ in the context of differential geometry was obtained in [20].

Now we are going to introduce a finer growth characteristics: multivariate growth and multivariate cogrowth.

Let $\Sigma = \{a_1, \dots, a_d\}$, $d \geq 2$ be an alphabet, Σ^* be the set of all finite words (or strings) over Σ . The set Σ^* with concatenation as a binary operation, can be interpreted as a monoid (with empty word serving as the identity element). In fact, Σ^* is a free monoid. Its growth sequence is d^n , $n = 0, 1, \dots$. Any subset $L \subset \Sigma^*$ is called a (formal) language. With any $w \in \Sigma^*$ we can associate the length $|w|$ and the frequency vector $\wp(w) = (|w|_{a_1}, \dots, |w|_{a_d}) \in \mathbb{N}^d$ where $|w|_{a_i}$ is a number of occurrences of the symbol a_i in the word w . Let $\mathbf{z} = (z_1, \dots, z_d) \in \mathbb{C}^d$ and

$$\Gamma_L(\mathbf{z}) = \sum_{w \in L} \mathbf{z}^{\wp(w)}, \quad (1)$$

where $\mathbf{z}^{\wp(w)} = z_1^{|w|_{a_1}} \dots z_d^{|w|_{a_d}}$. The series (1) is a multivariate series associated with L and can be rewritten as

$$\Gamma_L(\mathbf{z}) = \sum_{\mathbf{i} \in \mathbb{N}^d} \gamma_{\mathbf{i}} \mathbf{z}^{\mathbf{i}}, \quad (2)$$

where $\gamma_{\mathbf{i}}$ is the number of words w in L with the same vector $\wp(w) = \mathbf{i}$.

If we normalize the vector $\wp(w)$ as $\frac{1}{|w|} \wp(w)$ we get a vector of frequencies $\tilde{\wp}(w)$ belonging to the simplex M_d of probability vectors:

$$M_d = \left\{ \mathbf{r} \in \mathbb{R}_{\geq 0}^d : \mathbf{r} = (r_1, \dots, r_d), r_i \geq 0, \|\mathbf{r}\|_1 = \sum_{i=1}^d r_i = 1 \right\}$$

and $\|\cdot\|_1$ is a l_1 norm. The multivariate growth *indicatrice* that we are going to introduce is the number $\psi(\mathbf{r})$, $\mathbf{r} \in M_d$ which characterizes the growth of coefficients $\gamma_{\mathbf{i}}$ when $\|\mathbf{i}\|_1 \rightarrow \infty$ in the direction of vector \mathbf{r} . When $\mathbf{r} \in \mathbb{Q}^d$ is a rational vector, then the possible approach would be to define $\psi(\mathbf{r})$ by

$$\psi(\mathbf{r}) = \limsup_{n \rightarrow \infty} \frac{1}{n} \log |\gamma_{n\mathbf{r}}| \quad (3)$$

where $n \in \mathbb{N}$ and $\gamma_{n\mathbf{r}} = 1$ if $n\mathbf{r} \notin \mathbb{N}^d$. The Definition 4.1 given in the Section 4 follows the idea of [18]. It works for arbitrary $\mathbf{r} \in M_d$ and coincides with (3) in the rational case for many examples.

A crucial assumption that we make is the ‘‘concavity’’ assumption (CG) (see Definition 4.2) which allows to apply the powerful methods from convex analysis (as well as the results from [18]). In fact our definition works for arbitrary multivariate power series (2) with real coefficients. The main point is to present the indicatrice of growth $\psi(\mathbf{r})$ (assuming the condition (CG)) as

$$\psi(\mathbf{r}) = \inf_{\theta \in \partial(\Omega')} \langle \mathbf{r}, \theta \rangle \quad (4)$$

where $\Omega' = -\Omega \subset \mathbb{R}_{>0}^d$ and Ω is a closed convex set representing the Relog image of the domain of absolute convergence of (2) where

$$\text{Relog}(\mathbf{z}) = (\log |z_1|, \dots, \log |z_d|)$$

(see Theorem 4.1). We apply (4) for two languages: the Fibonacci language L_{Fib} and the language L_{F_m} of freely reduced words associated with a free group F_m of rank $m \geq 2$. These languages belong to class of regular languages, that is languages accepted by finite automaton. Regular languages play important role in many areas of mathematics, including dynamical systems and algebra. Regular normal form of elements in the group is a bijective presentation of elements of the group by elements of a regular language over the alphabet of generators and inverses. Regular geodesic normal form is such presentation for which the length of the element with respect to generating set coincides with the length of the word. Virtually abelian groups and Gromov hyperbolic groups have a regular geodesic normal form for any system of generators.

Regular languages are good in particular because their growth series (in one or multivariate case) are rational functions. This fact even in stronger form was known already to Chomsky and Schützenberger [3]. Proposition 2.1 of Section 2 gives a rational expression for a multivariate growth series associated with regular language. The condition (CG) mentioned earlier holds under the assumption of ergodicity of the automaton presenting the language (condition (E) in Section 2). It is satisfied in the presented examples and we summaries the computations from Section 5 and the Propositions 6.1 as

Theorem 1.1. *The indicatrices $\psi_{F_2}(\mathbf{r})$ and $\psi_{Fib}(\mathbf{r})$ are given by*

1.

$$\begin{aligned} \psi_{F_2}(\mathbf{r}) = \mathbf{H}(\mathbf{r}) + p \log \left(2q - p + 2\sqrt{p^2 - pq + q^2} \right) \\ + q \log \left(2p - q + 2\sqrt{p^2 - pq + q^2} \right), \end{aligned}$$

2.

$$\begin{aligned} \psi_{Fib}(\mathbf{r}) &= p \log \left(\frac{p}{p-q} \right) + q \log \left(\frac{p-q}{q} \right) \quad \text{if } p \geq \frac{1}{2}, \\ &= -\infty \quad \text{if } p < \frac{1}{2}, \end{aligned}$$

where $\mathbf{r} = (p, q) \in M_2$ and $\mathbf{H}(\mathbf{r}) = -p \log p - q \log q$ is the Shannon's entropy.

In fact, $\psi_{F_2}(\mathbf{r})$ is computed for modified multivariate growth series $\Delta_{F_2}(\mathbf{z})$, $\mathbf{r} \in M_2$ (see Section 3). The graphics of these functions are presented by Figure 2.

Regular languages quite often appear in Dynamical systems, first of all as languages associated with subshifts of finite type (SFT). For example a subshift of a full shift $\Sigma^{\mathbb{Z}}$ determined by a finite set \mathbf{F} of forbidden words; $\Sigma = \{0, 1\}$,

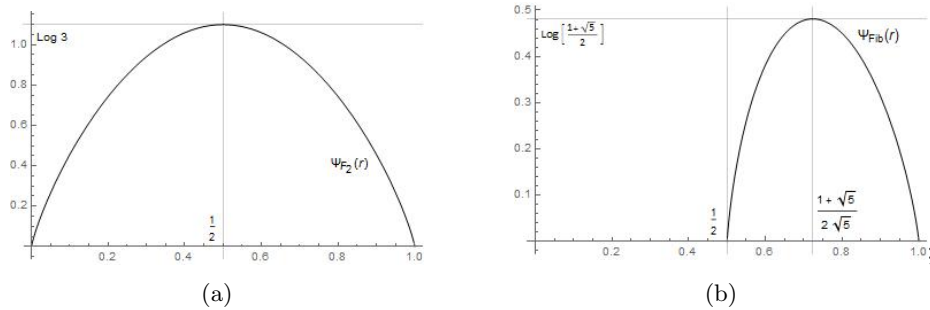


Figure 2: Graphs of $\psi_{F_2}(\mathbf{r})$ and $\psi_{Fib}(\mathbf{r})$.

$\mathbf{F} = \{11\}$, is the Fibonacci language. And if $\Sigma = \{a_1, \dots, a_m, a_1^{-1}, \dots, a_m^{-1}\}$, $\mathbf{F} = \{a_i a_i^{-1}, a_i^{-1} a_i : i = 1, \dots, m\}$ we get a subshift corresponds to the language L_{F_m} of freely reduced words representing elements of the free group F_m over Σ .

A standard way to present SFT is to define it by a digraph $\Gamma = (V, E)$ (i.e. a directed graph with allowed loops and multiple edges) or equivalently by a matrix A with non-negative integer entries. The theory of positive matrices based first of all on the powerful Perron-Frobenius theorem allows comprehensive study of SFT from the dynamical and other points of view. There is a canonical way to associate with irreducible SFT an ergodic Markov shift (and a Markov chain on Σ) given by a stochastic matrix P . Let (Δ, ζ) be a SFT where $\Delta \subset \Sigma^{\mathbb{Z}}$ and $\zeta : \Sigma^{\mathbb{Z}} \rightarrow \Sigma^{\mathbb{Z}}$ is a shift map defined by $(\zeta x)_n = x_{n+1}$, $x \in \Sigma^{\mathbb{Z}}$. Let $L(\Delta)$ be the language of (Δ, ζ) (all finite words that appears as a subwords in $x \in \Sigma^{\mathbb{Z}}$). Let $\psi_{\Delta}(\mathbf{r})$ be indicatrice of growth of $L(\Delta)$. The maximum of values of $\psi_{\Delta}(\mathbf{r})$ is a *topological entropy* of $L(\Delta)$ [15] and $\psi_{\Delta}(\mathbf{r})$ satisfies the relation (4), because $L(\Delta)$ is a regular language and the condition (CG) holds. There is an alternative way to present $\psi_{\Delta}(\mathbf{r})$ as the support function of a convex set via the methods of Large Deviation Theory (LDT) [5]. For our purpose it is enough only to apply Sanov's theorem stating that the Large Deviation Principle holds for finite Markov chains. The discussion in Section 7 relates formula (4) with the Sanov's expression for the rate function $I(\mathbf{r})$.

The paper is organized as follows. In Section 2, we recall some of the basic definitions from the theory of finite automata and formal languages that will be needed later. Then we introduce the condition (E) for the regular languages. The Section 3 is devoted to the computations of the modified multivariate growth series of language L_{F_m} of reduced elements of a free group F_m using two approaches. In Section 4, we discuss the condition (CG) and then prove Theorem 4.1. In Sections 5 and 6 we present computations of indicatrice $\psi(\mathbf{r})$ for F_2 and for the Fibonacci language L_{Fib} , respectively. Section 7 is devoted to application of Large Deviations Theory. Using Sanov Theorem, we get a relationship between $\psi(\mathbf{r})$ and the rate function $I(\mathbf{r})$. Using result from asymptotic

combinatorics in several variables presented in [16], in Section 8, we get a finer asymptotics associated with F_2 . Finally, Section 9 contains concluding remarks and some open questions.

2 Regular languages, their growth series and the condition (E)

We begin this section with recall of definition of finite automaton and language accepted by it. A finite automaton \mathcal{A} is given by a quintuple $(Q, \Sigma, \kappa, q_0, \mathcal{F})$, where Q is a finite set whose elements are called states, Σ is a finite alphabet, $\kappa : Q \times \Sigma \rightarrow Q$ is a transition function, the state $q_0 \in Q$ is a special state called the initial state and the set $\mathcal{F} \subset Q$ is nonempty set whose elements are called final states. It is convenient to visualize \mathcal{A} as a labeled directed graph $\Theta_{\mathcal{A}}$ with the vertex set Q , edge set

$$E = \{(q, s) : q, s \in Q, \kappa(q, a_i) = s \text{ for some } a_i \in \Sigma\},$$

and each such edge (q, s) with $\kappa(q, a_i) = s$ is supplied by the label a_i . Multiple edges and loops are allowed. The graph $\Theta_{\mathcal{A}}$ is called the diagram of \mathcal{A} . The example of these diagrams are presented by Figures 3a and 8. Observe that so defined \mathcal{A} is deterministic and complete automaton, i.e. given any $q \in Q$ and any $a \in \Sigma$ we know what would be the next state $\kappa(q, a)$. A word $w \in \Sigma^*$ is accepted by \mathcal{A} if starting with the initial state q_0 and traveling in diagram $\Theta_{\mathcal{A}}$ along with the path p_w determined by w we end up at some final state. Let $\mathcal{L}(\mathcal{A})$ be the set of words accepted by \mathcal{A} . A language $L \subset \Sigma^*$ is called regular if there is a finite automaton \mathcal{A} such that $L = \mathcal{L}(\mathcal{A})$. The important feature of this definition is the uniqueness of the path p_w for each $w \in \Sigma^*$.

One can generalize the above definition by replacing a singleton $\{q_0\}$ by a nonempty subset $\mathcal{I} \subset Q$ whose elements are called initial states and defining $\mathcal{L}(\mathcal{A})$ as a set of words w for which there is an initial state $i \in \mathcal{I}$ such that the path $p_{i,w}$ that begins at i and follow the word w ends up at \mathcal{F} . Surprisingly, this does not lead to a larger class of languages (as it is always possible to replace \mathcal{A} by automaton \mathcal{A}' with a single initial state such that $\mathcal{L}(\mathcal{A}) = \mathcal{L}(\mathcal{A}')$). The automaton with a single initial state are unambiguous in the sense that for each $w \in \mathcal{L}(\mathcal{A})$ there a unique path p_w that recognizes w . Nevertheless, there are situations (for instance in the case of the language of freely reduced words over the alphabet of generators of a free group) when ambiguous automata work better (see Figure 3b and Section 5).

One also can consider non-deterministic and incomplete automata. Still, this much larger class of automata determines the same class of languages, the class of regular languages. Non-deterministic automata appear for instance in the study of languages associated with sofic subshifts [15].

Recall that given $w \in \Sigma^*$ the vector $\wp(w) = (|w|_{a_1}, \dots, |w|_{a_d})$, where $|w|_{a_i}$ denote the number of occurrences of $a_i \in \Sigma$ in w . Let also $|w|$ denote the length

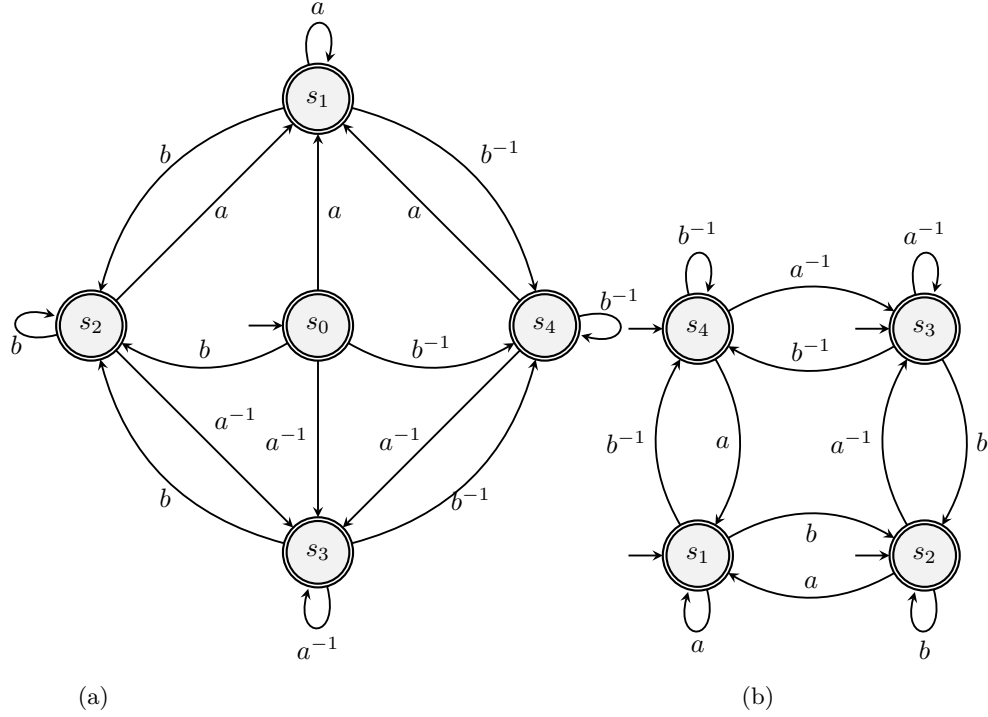


Figure 3: Diagrams of automata \mathcal{A} and \mathcal{A}' associated with the language L_{F_2} of reduced words representing elements of the free group F_2 , where the initial states are denoted by horizontal unlabeled arrow and the final states are represented by “double” circles.

of w . With $L \subset \Sigma^*$ one can associate a number of formal series: the ordinary series,

$$\Gamma_L(z) = \sum_{w \in L} z^{|w|}, \quad z \in \mathbb{C} \quad (5)$$

the multivariate series,

$$\Gamma_L(\mathbf{z}) = \Gamma_L(z_1, \dots, z_d) = \sum_{w \in L} \mathbf{z}^{\phi(w)}, \quad (6)$$

where $\mathbf{z} = (z_1, \dots, z_d) \in \mathbb{C}^d$ and $\mathbf{z}^{\phi(w)} = z_1^{|w|_{a_1}} \dots z_d^{|w|_{a_d}}$. Also one can consider the formal sum

$$\sum_{w \in L} w z^{|w|} \in \mathbb{Z}[w^*][[z]] \quad (7)$$

viewed as a formal power series with coefficients in the ring $\mathbb{Z}[w^*]$ (the semi-group ring of the free semi-group Σ^*). The consideration of these type of series go back at least to 50's of 20th century and is related first of all with the names

of Chomsky and Schützenberger [3].

For us it is important that in the case when L is a regular language the series (5), (6) and (7) are rational i.e. ratio of two polynomials. We focus our attention only on (5) and (6), and in fact the study of asymptotic properties of (6) is the main purpose of this article.

Let $A = (a_{ij})$ be the adjacency matrix of $\Theta_{\mathcal{A}}$ i.e. a $|Q| \times |Q|$ matrix whose rows and columns correspond to the states and

$$a_{ij} = \text{the number of edges in } \Gamma_{\mathcal{A}} \text{ joining } i \text{ with } j, \text{ where } i, j \in Q.$$

We use the ordering on Q in such a way that the first state is q_0 . Let

$$u = (1, 0, \dots, 0), \text{ and } v = (v_1, \dots, v_{|Q|})^t$$

be a row and column vectors of dimension $|Q|$, where

$$v_q = 1, \text{ if state } q \in \mathcal{F} \text{ and } v_q = 0 \text{ if } q \notin \mathcal{F}.$$

Then the standard technique of counting paths (see Theorem 4.7.2 of [19]) in finite graph (or in finite Markov chain) leads to the

$$\Gamma_L(z) = u \left[\sum_{n \geq 0} (zA)^n \right] v = u [I - zA]^{-1} v$$

and the rationality of $L(z)$ is obvious. A similar argument works for multivariate case, only the matrix zA should be replaced by

$$A(\mathbf{z}) = (a_{st}(\mathbf{z}))_{s,t=1}^{|Q|},$$

where $\mathbf{z} = (z_1, \dots, z_d)$ and

$$a_{st}(\mathbf{z}) = \begin{cases} \sum_i z_i & \text{where summation is taken over such } i \text{ that } \kappa(s, a_i) = t \\ 0 & \text{if there is no edge from } s \text{ to } t. \end{cases}$$

With this notations we have the following.

Proposition 2.1. *The multivariate growth series of a regular language $L = \mathcal{L}(\mathcal{A})$ satisfies*

$$\Gamma_L(\mathbf{z}) = u [I - A(\mathbf{z})]^{-1} v. \quad (8)$$

Hence $\Gamma_L(\mathbf{z}), \mathbf{z} \in \mathbb{C}^d$ is a multivariate rational function i.e.

$$\Gamma_L(\mathbf{z}) = \frac{G(\mathbf{z})}{H(\mathbf{z})},$$

where $G(\mathbf{z}), H(\mathbf{z}) \in \mathbb{C}[z_1, \dots, z_d]$. The polynomials $G(\mathbf{z}), H(\mathbf{z})$ obtained in (8), can be calculated, for instance, using the Cramer's rule.

For example, in the case of automaton presented by Figure 8,

$$A(\mathbf{z}) = \begin{pmatrix} z_1 & z_2 \\ z_1 & 0 \end{pmatrix} \text{ and } \Gamma_L(\mathbf{z}) = \frac{1 + z_2}{1 - z_1 - z_1 z_2}.$$

For us the following condition will be crucial.

Definition 2.1. (*Condition (E)*). We say that a language $L \subset \Sigma^*$ satisfy the condition (E) if there exists an integer $N \geq 0$ such that, for every $U, V \in L$, there exists $w \in \Sigma^*$, $|w| \leq N$ with $UwV \in L$.

The example of the regular languages satisfying the condition (E) come from ergodic (or primitive) automata \mathcal{A} i.e. automata in which any state can be connected to another state by a path in the diagram.

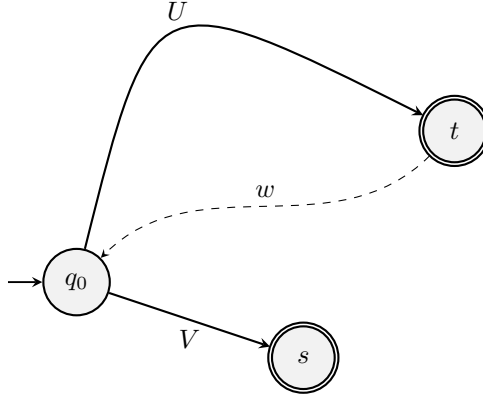


Figure 4: Part of Moore diagram of an ergodic automaton \mathcal{A}

This is because \mathcal{A} is ergodic automaton, $U, V \in \mathcal{L}(\mathcal{A})$ and s, t are end states of the path p_U, p_V then connecting state s with the initial state q_0 by some path q whose length does not exceed the diameter D of the graph $\Theta_{\mathcal{A}}$ (i.e. maximum of combinatorial distances between states in \mathcal{A}) we get $UwV \in \mathcal{L}(\mathcal{A})$ where w is a word read along path q , $|w| \leq D$. See Figure 4. So the condition (E) holds.

Unfortunately, not every regular language satisfying (E) can be accepted by ergodic automaton of the type described above. For instance the language L_{F_2} over the alphabet $\Sigma = \{a, a^{-1}, b, b^{-1}\}$ of freely reduced words (i.e. $aa^{-1}, a^{-1}a, bb^{-1}, b^{-1}b$ are forbidden) of the free group F_2 of rank 2, satisfy the condition (E) but can not be accepted by an ergodic unambiguous automaton [2]. The Figure 3 show two automata accepting the language L_{F_2} . The first automaton

presented in the Figure 3a is unambiguous but not ergodic. In the second automaton 3b, all states are initial, it is ergodic automaton but not unambiguous.

Later, we will introduce a condition (CG) for multivariate power series and it will follow that $\Gamma_L(\mathbf{z}), \mathbf{z} \in \mathbb{C}^d$ satisfies the condition (CG) if the associated language L satisfies the condition (E).

3 Computations of the multivariate growth series in the case of a free group

Let $F_m = \langle a_1, \dots, a_m \rangle$ be a free group of rank $m \geq 2$ and L_{F_m} be the language of freely reduced words over alphabet $\Sigma = \{a_1, \dots, a_m, a_1^{-1}, \dots, a_m^{-1}\}$. This is a regular language (an automaton accepting it in the case $m = 2$ is given by Figure 3a) and hence a corresponding multivariate growth series $\Gamma_{F_m}(\mathbf{z})$ is rational. It can be computed using for instance Proposition 2.1. Even in the case of $m = 2$ computations are quite involved. So instead we will focus on a modified version $\Delta_{F_m}(\mathbf{z})$ defined as

$$\Delta_{F_m}(\mathbf{z}) = 1 + \sum_{g \in F_m, g \neq e} \mathbf{z}^{\varphi(g)}$$

where $\varphi(g) = (|g|_{a_1^{\pm}}, \dots, |g|_{a_m^{\pm}})$ and $|g|_{a_i^{\pm}}$ is a number of occurrences of symbols a_i and a_i^{-1} in the freely reduced word presenting element g i.e. $|g|_{a_i^{\pm}} = |g|_{a_i} + |g|_{a_i^{-1}}$. The modified series $\Delta_{F_m}(\mathbf{z})$ is also important for us, because of the question 9.1 formulated at the end of the article in connection with symmetric random walks. Of course if we know $\Gamma_{F_m}(\mathbf{z})$, then we can obtain the $\Delta_{F_m}(\mathbf{z})$ and vice versa.

We present the computations of the modified multivariate growth series $\Delta_{F_m}(\mathbf{z})$, first using the automata approach i.e using the Proposition 2.1 and later, using the symmetries of this language leading to a simpler system of equations and as a result of simpler expression for $\Delta_{F_m}(\mathbf{z})$ given by Proposition 3.2.

Proposition 3.1.

$$\Gamma_{F_m}(\mathbf{z}) = \frac{(1 + z_1) \cdots (1 + z_m)}{R(\mathbf{z})} \quad (9)$$

where

$$R(\mathbf{z}) = 1 - \left(\sum_i z_i + 3 \sum_{i < j} z_i z_j + \cdots + (2l - 1) \sum_{i_1 < \cdots < i_l} \prod_{k=1}^l z_{i_k} + \cdots + (2m - 3) \sum_{i_1 < \cdots < i_{m-1}} \prod_{k=1}^{m-1} z_{i_k} + (2m - 1) z_1 \cdots z_m \right).$$

Proof. Let \mathcal{A}_{F_m} be an automaton generating language L_{F_m} of reduced elements of F_m . See Figure 3a for \mathcal{A}_{F_2} . The $2m+1$ states of \mathcal{A}_{F_m} are $s_0, s_1, \dots, s_m, s_{m+1} = s_1^{-1}, \dots, s_{2m} = s_m^{-1}$ and the adjacency matrix A of \mathcal{A}_{F_m} is

$$A = \begin{pmatrix} \mathbf{0}_{1 \times 1} & J_{1 \times m} & J_{1 \times m} \\ \mathbf{0}_{m \times 1} & J_{m \times m} & (J - I)_{m \times m} \\ \mathbf{0}_{m \times 1} & (J - I)_{m \times m} & J_{m \times m} \end{pmatrix},$$

where $\mathbf{0}_{m \times n}$ is zero matrix, $J_{m \times n}$ is matrix with all entries are one and $(J - I)_{m \times m}$ is matrix having diagonal entries are zero and one elsewhere. Using the above decomposition of A , we obtain the transfer matrix $A(\mathbf{z})$ and then applying Equation (8) of the Proposition 2.1 we obtain

$$\Gamma_{F_m}(\mathbf{z}) = \frac{\sum_i (-1)^{i+1} \det(I - A(\mathbf{z}) : i, 1)}{\det(I - A(\mathbf{z}))}, \quad (10)$$

where sum is taken over all final states s_i and if N is matrix then $(N : i, j)$ denotes the matrix obtained by removing the i th row and j th column of N . The matrix I is the $2m + 1$ dimensional identity matrix.

We divide the rest of the proof in two lemmas.

Lemma 3.1.

$$\det(I - A(\mathbf{z})) = (1 - z_1) \cdots (1 - z_m) \times R(\mathbf{z})$$

Proof. We prove the lemma using induction on $m \geq 2$. It is easy to see that the statement is true when $m = 2$. i.e.

$$\det(I - A(\mathbf{z})) = (1 - z_1)(1 - z_2)(1 - z_1 - z_2 - 3z_1z_2).$$

Assume that the statement is true when $m = k$. i.e.

$$\begin{aligned} \det(I - A(\mathbf{z})) &= (1 - z_1) \cdots (1 - z_k) \times \\ &\left(1 - \left(\sum_i z_i + 3 \sum_{i < j} z_i z_j + \cdots + (2l - 1) \sum_{i_1 < \cdots < i_l} \prod_{j=1}^l z_{i_j} \right. \right. \\ &\left. \left. + \cdots + (2k - 3) \sum_{i_1 < \cdots < i_{k-1}} \prod_{j=1}^{k-1} z_{i_j} + (2k - 1) z_1 \cdots z_k \right) \right). \end{aligned}$$

Let $m = k + 1$. Then $I - A(\mathbf{z})$

$$= \begin{pmatrix} 1 & -z_1 & \cdots & -z_{k+1} & -z_1 & \cdots & -z_{k+1} \\ 0 & 1 - z_1 & \cdots & -z_{k+1} & 0 & \cdots & -z_{k+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & -z_1 & \cdots & 1 - z_{k+1} & -z_1 & \cdots & 0 \\ 0 & 0 & \cdots & -z_{k+1} & 1 - z_1 & \cdots & -z_{k+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & -z_1 & \cdots & 0 & -z_1 & \cdots & 1 - z_{k+1} \end{pmatrix}$$

We rearrange the rows and columns of $I - A(\mathbf{z})$ by interchanging the positions of $(k + 1)$ th row and column with $\{k + 2, \dots, 2k + 1\}$ th rows and columns, respectively. Then $I - A(\mathbf{z})$

$$= \left(\begin{array}{c|cccccc|cc} 1 & -z_1 & \cdots & -z_k & -z_1 & \cdots & -z_k & -z_{k+1} & -z_{k+1} \\ \hline 0 & 1 - z_1 & \cdots & -z_k & 0 & \cdots & -z_k & -z_{k+1} & -z_{k+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & -z_1 & \cdots & 1 - z_k & -z_1 & \cdots & 0 & -z_{k+1} & -z_{k+1} \\ 0 & 0 & \cdots & -z_k & 1 - z_1 & \cdots & -z_k & -z_{k+1} & -z_{k+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & -z_1 & \cdots & 0 & -z_1 & \cdots & 1 - z_k & -z_{k+1} & -z_{k+1} \\ \hline 0 & -z_1 & \cdots & -z_k & -z_1 & \cdots & -z_k & 1 - z_{k+1} & 0 \\ 0 & -z_1 & \cdots & -z_k & -z_1 & \cdots & -z_k & 0 & 1 - z_{k+1} \end{array} \right)$$

Observe that the rows and columns of $I - A(\mathbf{z})$ are indexed as $s_0, s_1, \dots, s_k, s_1^{-1}, \dots, s_k^{-1}, s_{k+1}, s_{k+1}^{-1}$ and the determinant of the central block and the expression that we have assumed when $m = k$ are identical. As first column has exactly one 1 and zero elsewhere, we can ignore the first row and first column. Hence $\det(I - A(\mathbf{z}))$

$$= \det \left(\begin{array}{c|cccccc|cc} 1 - z_1 & \cdots & -z_k & 0 & \cdots & -z_k & -z_{k+1} & -z_{k+1} \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ -z_1 & \cdots & 1 - z_k & -z_1 & \cdots & 0 & -z_{k+1} & -z_{k+1} \\ 0 & \cdots & -z_k & 1 - z_1 & \cdots & -z_k & -z_{k+1} & -z_{k+1} \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ -z_1 & \cdots & 0 & -z_1 & \cdots & 1 - z_k & -z_{k+1} & -z_{k+1} \\ \hline -z_1 & \cdots & -z_k & -z_1 & \cdots & -z_k & 1 - z_{k+1} & 0 \\ -z_1 & \cdots & -z_k & -z_1 & \cdots & -z_k & 0 & 1 - z_{k+1} \end{array} \right)$$

We rewrite the matrix by naming these blocks as

$$\det(I - A(\mathbf{z})) = \det \begin{pmatrix} A_{2k \times 2k} & B_{2k \times 2} \\ C_{2 \times 2k} & D_{2 \times 2} \end{pmatrix}$$

Recall that $z_i \neq 1$, for all $i = 1, \dots, k + 1$ and hence $\det(D_{2 \times 2}) \neq 0$. Therefore we have

$$\det(I - A(\mathbf{z})) = \det \begin{pmatrix} A'_{2k \times 2k} & B_{2k \times 2} \\ \mathbf{0}_{2 \times 2k} & D_{2 \times 2} \end{pmatrix} = \det(A'_{2k \times 2k}) \det(D_{2 \times 2}),$$

where $A' = A_{2k \times 2k} - B_{k \times 2} D_{2 \times 2}^{-1} C_{2 \times k}$ and its (i, j) th entry

$$A'(i, j) = \begin{cases} 1 - z_i - \frac{2z_i z_{k+1}}{1 - z_{k+1}} & \text{if } i = j \\ -\frac{2z_i z_{k+1}}{1 - z_{k+1}} & \text{if } i < j \text{ and } j = i + k(\text{mod } 2k) \\ -\frac{2z_i z_{k+1}}{1 - z_{k+1}} & \text{if } j < i \text{ and } i = j + k(\text{mod } 2k) \\ -z_i - \frac{2z_i z_{k+1}}{1 - z_{k+1}} & \text{otherwise.} \end{cases}$$

In order to prove the $m = k + 1$ step, it suffices to show that

$$\det(A') = \frac{(1 - z_1) \cdots (1 - z_k)}{1 - z_{k+1}} \times \left(1 - \sum_i z_i - 3 \sum_{i < j} z_i z_j - \cdots - (2k - 1) \sum_{i_1 < \cdots < i_k} \prod_{j=1}^k z_{i_j} - (2k + 1) z_1 \cdots z_{k+1} \right).$$

Applying below rows and columns transformations

1. $R_i - R_1, i = 2, \dots, 2k$
2. $C_i - C_{k+i}, i = 1, \dots, k$
3. $R_i + R_1, i = 2, \dots, k, k + 2, \dots, 2k$ and $R_{k+1} + 2R_1$
4. $R_{k+i} + R_i, i = 2, \dots, k$
5. $R_i + R_1, i = 2, \dots, k$ and $R_i + 2R_1, i = k + 1, \dots, 2k$

to the matrix A' , we convert A' to the below upper triangular block matrix.

$$\det(A') = \det \left(\begin{array}{ccc|ccc} 1 - z_1 & \cdots & 0 & -\frac{2z_1 z_{k+1}}{1 - z_{k+1}} & \cdots & -z_1 - \frac{2z_k z_{k+1}}{1 - z_{k+1}} \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & \cdots & 1 - z_k & -z_k - \frac{2z_1 z_{k+1}}{1 - z_{k+1}} & \cdots & -\frac{2z_k z_{k+1}}{1 - z_{k+1}} \\ \hline 0 & \cdots & 0 & 1 - z_1 - \frac{4z_1 z_{k+1}}{1 - z_{k+1}} & \cdots & 2 \left(-z_k - \frac{2z_k z_{k+1}}{1 - z_{k+1}} \right) \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & 2 \left(-z_1 - \frac{2z_1 z_{k+1}}{1 - z_{k+1}} \right) & \cdots & 1 - z_k - \frac{4z_k z_{k+1}}{1 - z_{k+1}} \end{array} \right)$$

We call the bottom right corner block as A'' . In A'' matrix, subtracting twice the first row from each row $R_i, i = 2, \dots, k$, we get

$$\det(A'') = \det \left(\begin{array}{cccc} \frac{1 - z_1 - z_{k+1} - 3z_1 z_{k+1}}{1 - z_{k+1}} & \frac{-2z_2(1 + z_{k+1})}{1 - z_{k+1}} & \cdots & \frac{-2z_k(1 + z_{k+1})}{1 - z_{k+1}} \\ -1 - z_1 & 1 + z_2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ -1 - z_1 & 0 & \cdots & 1 + z_k \end{array} \right)$$

$$= \left(1 - \sum_i z_i - 3 \sum_{i < j} z_i z_j - \cdots - (2k-1) \sum_{i_1 < \cdots < i_k} \prod_{j=1}^k z_{i_j} - (2k+1) z_1 \cdots z_{k+1} \right).$$

This implies that the statement is true for $m = k + 1$ and hence the lemma is proved. \square

Lemma 3.2.

$$\sum_i (-1)^{i+1} \det(I - A(\mathbf{z}) : i, 1) = (1 - z_1^2) \cdots (1 - z_m^2).$$

Proof. Using similar strategy as in the above proof, we can show that

$$\begin{aligned} \det(I - A(\mathbf{z}) : i + 1, 1) &= \det(I - A(\mathbf{z}) : i + m + 1, 1) \\ &= z_i (1 - z_i) \prod_{j \neq i} (1 - z_j^2) \end{aligned}$$

This implies

$$\sum_i (-1)^{i+1} \det(I - A(\mathbf{z}) : i, 1) = (1 - z_1^2) \cdots (1 - z_m^2).$$

\square

Hence, the proof of Proposition 3.1 follows from Lemmas 3.1 and 3.2. \square

We now provide an alternative way of computations.

Proposition 3.2. *The function $\Delta_{F_m}(\mathbf{z})$ satisfies*

$$\Delta_{F_m}(\mathbf{z}) = \frac{1}{1 - 2 \sum_{i=1}^m \frac{z_i}{1 + z_i}}. \quad (11)$$

Proof. Recall that the elements of F_m are identified with freely reduced words over the alphabet $\Sigma = \{a_1, \dots, a_m, a_1^{-1}, \dots, a_m^{-1}\}$.

$$\Delta_{F_m}(\mathbf{z}) = 1 + \sum_{e \neq g \in F_m} \mathbf{z}^{\wp(g)}$$

where e is the identity in F_m , $\mathbf{z} = (z_1, \dots, z_m)$, $\wp(g) = (|g|_{a_1^\epsilon}, \dots, |g|_{a_m^\epsilon})$ and $\mathbf{z}^{\wp(g)} = z_1^{|g|_{a_1^\epsilon}} \cdots z_m^{|g|_{a_m^\epsilon}}$.

For each $i = 1, \dots, m$, $\epsilon = \pm 1$, let $F_{m,i}^\epsilon = \{g \in F_m \setminus \{e\} : \text{the first letter of } g \text{ is } a_i^\epsilon\}$. Define

$$\Delta_i^\epsilon(\mathbf{z}) = \Delta_{F_{m,i}^\epsilon}^\epsilon(\mathbf{z}) = \sum_{e \neq g \in F_{m,i}^\epsilon} \mathbf{z}^{\wp(g)}$$

so that we get

$$\Delta_{F_m}(\mathbf{z}) = 1 + \sum_{i,\epsilon} \Delta_i^\epsilon(\mathbf{z}).$$

Observe that

$$\Delta_i(\mathbf{z}) = z_i + z_i \Delta_i(\mathbf{z}) + z_i \sum_{j \neq i, \epsilon} \Delta_j^\epsilon(\mathbf{z}).$$

This implies

$$\Delta_i(\mathbf{z}) = z_i (\Delta_{F_m}(\mathbf{z}) - \Delta_i^{-1}(\mathbf{z})).$$

Similarly we can write

$$\Delta_i^{-1}(\mathbf{z}) = z_i (\Delta_{F_m}(\mathbf{z}) - \Delta_i(\mathbf{z})).$$

Adding $\Delta_i(\mathbf{z})$ and $\Delta_i^{-1}(\mathbf{z})$ we get

$$\Delta_i(\mathbf{z}) + \Delta_i^{-1}(\mathbf{z}) = \left(\frac{2z_i}{1+z_i} \right) \Delta_{F_m}(\mathbf{z}),$$

Hence

$$\Delta_{F_m}(\mathbf{z}) = 1 + \sum_i (\Delta_i(\mathbf{z}) + \Delta_i^{-1}(\mathbf{z})) = 1 + 2 \sum_{i=1}^m \left(\frac{z_i}{1+z_i} \right) \Delta_{F_m}(\mathbf{z})$$

and we come to (11). \square

We now compare two expressions for $\Delta_{F_m}(\mathbf{z})$. Multiplying numerator and denominator of (11) by $(1+z_1) \cdots (1+z_m)$ we get

$$\Delta_{F_m}(\mathbf{z}) = \frac{(1+z_1) \cdots (1+z_m)}{((1+z_1) \cdots (1+z_m)) - 2((1+z_1) \cdots (1+z_m)) \sum_{i=1}^m \frac{z_i}{1+z_i}} \quad (12)$$

It is easy to see that

$$(1+z_1) \cdots (1+z_m) = 1 + \sum_i z_i + \sum_{i_2 < i_1} z_{i_1} z_{i_2} + \cdots + \sum_{i_1 < \cdots < i_l} \prod_{k=1}^l z_{i_k} + \cdots + \prod_{i=1}^m z_i$$

and

$$\begin{aligned} & 2((1+z_1) \cdots (1+z_m)) \sum_{i=1}^m \frac{z_i}{1+z_i} \\ &= 2 \sum_i z_i + 4 \sum_{i_2 < i_1} z_{i_1} z_{i_2} + \cdots + 2l \sum_{i_1 < \cdots < i_l} \prod_{k=1}^l z_{i_k} + \cdots + 2m \prod_{i=1}^m z_i \end{aligned}$$

Substituting above relations in the denominator of (12) we get $R(\mathbf{z})$ of (9).

4 The multivariate growth exponent and the condition (CG)

Let

$$\Gamma(\mathbf{z}) = \sum_{\mathbf{i} \in \mathbb{N}^d} f_{\mathbf{i}} \mathbf{z}^{\mathbf{i}} = \sum_{i_1, \dots, i_d} f_{i_1, \dots, i_d} z_1^{i_1} \cdots z_d^{i_d} \quad (13)$$

be a multivariate power series. We denote by \mathcal{D} the interior of the domain of absolute convergence of Γ , which we assume to be non empty. In particular, 0 belongs to \mathcal{D} .

We are going to define for each $\mathbf{r} \in M_d$ a growth exponent $\varphi(\mathbf{r})$ for coefficients $\{f_{\mathbf{i}}\}_{\mathbf{i} \in \mathbb{N}^d}$ in the direction of \mathbf{r} . When \mathbf{r} is a vector with rational entries, then one can define $\varphi(\mathbf{r})$ as

$$\varphi(\mathbf{r}) = \limsup_{n \rightarrow \infty} |f_{n\mathbf{r}}|^{\frac{1}{n}}, \quad (14)$$

where only coefficients $n\mathbf{r} = (nr_1, \dots, nr_d) \in \mathbb{N}^d$ with integer entries are taken into account. This approach is considered in the book [16] and some of its references.

To define $\varphi(\mathbf{r})$ for arbitrary $\mathbf{r} \in M_d$ following [18], we first define the function $\psi(\mathbf{r})$. We then hope to be able put $\varphi(\mathbf{r}) = e^{\psi(\mathbf{r})}$, which would then agree with (3).

Definition 4.1. *Let $C \subset \mathbb{R}^d$ be an open cone at $\mathbf{0} \in \mathbb{R}^d$ containing vector $\mathbf{r} \in \mathbb{R}^d$. Define*

$$\tau_C = \limsup_{R \rightarrow \infty} \frac{1}{R} \log \left(\sum_{\substack{\mathbf{i} \in C \\ R \leq \|\mathbf{i}\|_1 \leq R+1}} f_{\mathbf{i}} \right) \quad (15)$$

and

$$\psi(\mathbf{r}) = \|\mathbf{r}\|_1 \inf_{\mathbf{r} \in C} \tau_C \quad (16)$$

where inf is taken over open cones containing \mathbf{r} .

We call $\psi(\mathbf{r})$ *indicatrice* of growth. The assumption that 0 is an interior point of the domain of absolute convergence of Γ ensures that ψ can not take the value $+\infty$, but it may take the value $-\infty$.

Recall that the ReLog map is a map $\text{ReLog} : \mathbb{C}_*^d \rightarrow \mathbb{R}^d$, where $\mathbb{C}_* = \mathbb{C} \setminus \{0\}$ given by

$$\text{ReLog}(\mathbf{z}) = (\log |z_1|, \dots, \log |z_d|).$$

In future, we will use the notation $\mathcal{E} = \mathbb{R}^d$ and \mathcal{E}^* for its dual space i.e. the space of continuous linear functionals $\theta : \mathcal{E} \rightarrow \mathbb{R}$ with a natural identification of

\mathcal{E}^* with \mathcal{E} via the standard inner product

$$\langle \mathbf{x}, \theta \rangle = \theta(\mathbf{x}) = \sum_{i=1}^d x_i \theta_i,$$

where $\mathbf{x} = (x_1, \dots, x_d) \in \mathcal{E}$, $\theta = (\theta_1, \dots, \theta_d) \in \mathcal{E}^*$.

The height function $h_{\mathbf{r}}(\mathbf{z})$ for $\mathbf{z} \in \mathcal{E}$ is the function

$$h_{\mathbf{r}}(\mathbf{z}) = - \sum_{i=1}^d r_i \log |z_i| = -\langle \mathbf{r}, \text{Re log}(\mathbf{z}) \rangle$$

if $\mathbf{r} \in M_d \cap \mathbb{Q}^d$ is a rational vector then for (14) there is an upper bound

$$\varphi(\mathbf{r}) = \limsup_{n \rightarrow \infty} |f_{n\mathbf{r}}|^{\frac{1}{n}} \leq |z'_1|^{-r_1} \dots |z'_d|^{-r_d} = e^{-h_{\mathbf{r}}(\mathbf{z}')} \quad (17)$$

which hold for every point \mathbf{z}' in the closure of $\overline{\mathcal{D}}$ of \mathcal{D} , as shown for instance in ([16], formula (5.15)). We shall provide sufficient conditions for replacement of the inequality in (17) by the equality for properly chosen \mathbf{z}' and extend the definition of $\varphi(\mathbf{r})$ to irrational \mathbf{r} by $\varphi(\mathbf{r}) = e^{\psi(\mathbf{r})}$. The function $\varphi(\mathbf{r})$ will be called the multivariate growth exponent.

For what follow it will be convenient to associate with the power series $\Gamma(\mathbf{z})$ a Radon measure $\nu = \nu_{\Gamma}$ defined on Borel subsets S of \mathcal{E} by

$$\nu(S) = \sum_{\mathbf{i} \in S} f_{\mathbf{i}} \quad (18)$$

We shall now define the condition (CG).

Definition 4.2. (*Condition (CG)*) Following [18] we say that $\Gamma(\mathbf{z})$ has a concave growth of coefficients if there are $a, b, c > 0$ such that for all $\mathbf{x}, \mathbf{y} \in \mathcal{E}$

$$\nu(B_{\mathbf{x}+\mathbf{y}}(a)) \geq c\nu(B_{\mathbf{x}}(b))\nu(B_{\mathbf{y}}(b)) \quad (19)$$

where $B_{\mathbf{x}}(a)$ is the ball of radius a with the center at $\mathbf{x} \in \mathcal{E}$ for the norm $\|\cdot\|_1$.

A large class of power series $\Gamma(\mathbf{z})$ satisfying the condition (CG) are generating series of regular languages satisfying the condition (E), as explained in section 2.

Given a Radon measure ν satisfying condition (CG) one defines for a given open cone $C \subset \mathcal{E}$ with a root at $\mathbf{0}$,

$$\tau_{\nu, C} = \limsup_{R \rightarrow \infty} \log \nu(B_{\mathbf{0}}(R+1) \setminus B_{\mathbf{0}}(R)),$$

$$\tau_{\nu}(\mathbf{r}) = \|\mathbf{r}\|_1 \inf_{\mathbf{r} \in C} \tau_{\nu, C}$$

for $\mathbf{r} \in \mathcal{E}$ and $\|\mathbf{r}\|_1 = 1$ and finally for $\mathbf{x} \in \mathcal{E}$

$$\psi(\mathbf{x}) = \begin{cases} \tau_\nu \left(\frac{\mathbf{x}}{\|\mathbf{x}\|_1} \right) & \text{if } \mathbf{x} \neq \mathbf{0} \\ 0 & \text{if } \mathbf{x} = \mathbf{0}. \end{cases} \quad (20)$$

Then assuming that $\tau_\nu < \infty$, where

$$\tau_\nu = \sup_{\mathbf{x} \in \mathcal{E}, \|\mathbf{x}\|_1=1} \psi_\nu(\mathbf{x}) \quad (21)$$

we know from ([18], Lemma 3.1.7) that the function $\psi_\nu : \mathcal{E} \rightarrow \mathbb{R} \cup \{-\infty\}$ is upper semi-continuous. Obviously it is positively homogeneous i.e. for any $t > 0$

$$\psi_\nu(t\mathbf{x}) = t\psi_\nu(\mathbf{x}).$$

Finally, if the Condition (CG) holds, then $\psi_\nu(\mathbf{x})$ is concave i.e.

$$\psi_\nu(\mathbf{x} + \mathbf{y}) \geq \psi_\nu(\mathbf{x}) + \psi_\nu(\mathbf{y}), \text{ for } \mathbf{x}, \mathbf{y} \in \mathcal{E}.$$

Observe that in our situation, because the measure $\nu = \nu_\Gamma$ is supported only on the $\mathbb{N}^d \subset \mathbb{R}_{\geq 0}^d$ part of \mathcal{E} ,

$$\psi_\nu(\mathbf{x}) = -\infty, \text{ for } \mathbf{x} \notin \mathbb{R}_{\geq 0}^d.$$

Let $\Omega = \text{ReLog}(\mathcal{D} \cap \mathbb{C}_*^d) \subset \mathcal{E}$ (recall that $\mathcal{D} \subset \mathbb{C}^d$ is the interior of the domain of absolute convergence of (13)). It is well known ([16], Proposition 3.4) that Ω is convex. Let $\psi = \psi_\Gamma$ be the function given by Definition 4.1.

Theorem 4.1. *Let $\Gamma(\mathbf{z})$ be a power series given by the Equation (13) with non-negative coefficients $f_{\mathbf{i}}, \mathbf{i} \in \mathbb{N}^d$ such that $f_{\mathbf{i}}$ satisfy the concavity condition (CG) and $\tau_\nu < \infty$ where $\nu = \nu_\Gamma$ is given by the Equation (18) and τ_ν is given by (21). Let $\mathcal{D} \subset \mathbb{C}^d$ be the interior of the set of points of absolute convergence of $\Gamma(\mathbf{z})$, $\Omega = \text{ReLog}(\mathcal{D} \cap \mathbb{C}_*^d)$ and $\bar{\Omega}$ be the closure of Ω . Then $-\psi_\Gamma(\mathbf{x})$ is the support function of the closure $\bar{\Omega}$ and for $\mathbf{x} \in \mathbb{R}_{\geq 0}^d$*

$$\psi_\Gamma(\mathbf{x}) = \inf_{\theta \in -\bar{\Omega}} \langle \mathbf{x}, \theta \rangle \quad (22)$$

$$= \inf_{\theta \in -\partial\bar{\Omega}} \langle \mathbf{x}, \theta \rangle \text{ if } \Omega \neq \mathbb{R}^d. \quad (23)$$

Proof. Because of the Condition (CG), the function $-\psi_\Gamma(\mathbf{x})$ is a lower semi-continuous, convex and positively homogeneous. Hence, it is a support function of a one and only one closed convex subset $S \subset \mathcal{E}^*$ given by

$$\begin{aligned} S &= \{ \theta \in \mathcal{E}^* : -\psi_\Gamma(\mathbf{x}) \geq \langle \mathbf{x}, \theta \rangle, \forall \mathbf{x} \in \mathcal{E} \} \\ &= \{ \theta \in \mathcal{E}^* : \psi_\Gamma(\mathbf{x}) \leq -\langle \mathbf{x}, \theta \rangle, \forall \mathbf{x} \in \mathcal{E} \} \end{aligned} \quad (24)$$

The set of absolute convergence of $\Gamma(\mathbf{z})$ is determined by the condition that the integral given below in the Equation (25) is convergent.

$$\begin{aligned}\Gamma(|z_1|, \dots, |z_d|) &= \sum_{\mathbf{i}} f_{\mathbf{i}} e^{\sum_{k=1}^d i_k \log |z_k|} \\ &= \int_{\mathcal{E}} e^{\langle \mathbf{x}, \theta \rangle} d\nu_{\Gamma}(\mathbf{x})\end{aligned}\quad (25)$$

where $\theta = \text{Re log}(\mathbf{z}), \mathbf{z} = (z_1, \dots, z_d)$.

Consider the set

$$\Delta_{\Gamma}^{\circ} = \{\theta \in \mathcal{E}^* : -\langle \mathbf{x}, \theta \rangle > \psi_{\Gamma}(\mathbf{x}), \forall \mathbf{x} \in \mathcal{E} \setminus \{\mathbf{0}\}\}.$$

Using ([18], Lemma 3.1.3), we conclude that if $\theta \in \Delta_{\Gamma}^{\circ}$, then the integral (25) converges while if there is $\mathbf{x} \in \mathcal{E} \setminus \{\mathbf{0}\}$ with $-\langle \mathbf{x}, \theta \rangle < \psi_{\Gamma}(\mathbf{x})$, then it diverges. The closure $\Delta_{\Gamma} = \overline{\Delta_{\Gamma}^{\circ}}$ is the set

$$\Delta_{\Gamma} = \{\theta \in \mathcal{E}^* : -\langle \mathbf{x}, \theta \rangle \geq \psi_{\Gamma}(\mathbf{x}), \forall \mathbf{x} \in \mathcal{E}\}$$

which is the set S given by the Equation (24) for whom $-\psi_{\Gamma}(\mathbf{x})$ is the support function. Thus

$$\begin{aligned}-\psi_{\Gamma}(\mathbf{x}) &= \sup_{\theta \in S} \langle \mathbf{x}, \theta \rangle \\ \psi_{\Gamma}(\mathbf{x}) &= -\sup_{\theta \in S} \langle \mathbf{x}, \theta \rangle \\ &= \inf_{\theta \in -S} \langle \mathbf{x}, \theta \rangle\end{aligned}$$

Also, we conclude that

$$S = \text{Re log}(\overline{\mathcal{D}}) = \overline{\Omega}.$$

This leads us to the Equation (22) and eventually to (23) as the maximum of a linear functional taken over the closed convex set S is achieved on the boundary ∂S of S or equals to $-\infty$. \square

Following is the immediate consequence of Theorem 4.1.

Corollary 4.1.

$$\psi(\mathbf{x}) = \inf_{\mathbf{z} \in \partial \mathcal{D}} \left(-\sum x_i \log |z_i| \right) \quad (26)$$

Example

Let $d = 2, \Sigma = \{a_1, a_2\}$ and $L = \Sigma^*$ be the language of a free monoid. Then

$$\Gamma_L(\mathbf{z}) = \frac{1}{1 - z_1 - z_2} = \sum_{n=0}^{\infty} \sum_{i=0}^n \binom{n}{i} z_1^i z_2^{n-i}.$$

It is obvious that the Condition (CG) holds and the direct computation based on the use of Stirling's formula or Theorem 4.1 show that for $\mathbf{r} \in M_2$ the function $\psi_L(\mathbf{x})$ is the Shannon's entropy $\mathbf{H}(\mathbf{r})$. i.e.

$$\psi_L(\mathbf{r}) = \mathbf{H}(\mathbf{r}) = -r_1 \log r_1 - r_2 \log r_2, \mathbf{r} \in M_2.$$

More examples will be discussed in the following two sections.

5 Multivariate growth exponent in the case of the free group

In the section 3, we got two expressions for the multivariate growth series $\Gamma_{F_m}(\mathbf{z})$ for the language L_{F_m} of freely reduced words in the alphabet $\{a_1, \dots, a_m, a_1^{-1}, \dots, a_m^{-1}\}$.

Consider the case $m = 2$. Then

$$\Delta_{F_2}(\mathbf{z}) = \frac{(1+z_1)(1+z_2)}{1-z_1-z_2-3z_1z_2} \quad (27)$$

We use the notations $\mathcal{D}, \Omega, \text{Relog}$ of Section 4. First, we are mainly interested in understanding the shape of the set $\Omega = \text{Relog}(\mathcal{D})$. To get a view of the real slice of the domain of the absolute convergence of the power series (27), we observe that the domain \mathcal{D} is described as

$$\mathcal{D} = \{(z_1, z_2) \in \mathbb{C}^2 : |z_1| + |z_2| + 3|z_1||z_2| < 1\}$$

The real slice of the curve

$$H(z_1, z_2) = f_1(\mathbf{z}) = 1 - z_1 - z_2 - 3z_1z_2 = 0$$

is presented by hyperbola in Figure 5a.

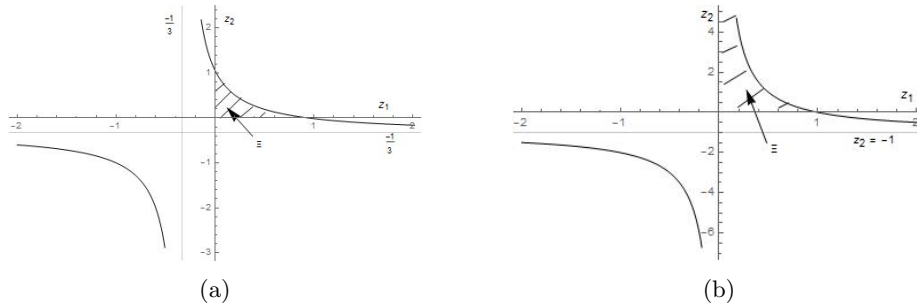


Figure 5: Real part of the Domain \mathcal{D} of $H(z_1, z_2)$ of L_{F_2} and L_{Fib} , respectively.

The real slice of \mathcal{D} with $z_1, z_2 \in \mathbb{R}, z_1, z_2 \geq 0$ is presented by the tented area Ξ . The set $-\Omega$ is obtained from Ξ by making the substitution $z_1 = e^{-s}, z_2 = e^{-t}$. To get the clear picture of the shape of the set Ω (and hence $-\Omega$) we make use of an interesting notion from algebraic geometry called *amoeba*.

Recall, that given a Laurent polynomial $f(\mathbf{z}), \mathbf{z} \in \mathbb{C}^d$, the amoeba of f is the set

$$\text{amoeba}(f) = \{\text{ReLog}(\mathbf{z}) : \mathbf{z} \in \mathbb{C}_*^d, f(\mathbf{z}) = 0\} \subset \mathbb{R}^d,$$

where $\mathbb{C}_* = \mathbb{C} \setminus \{0\}$. The amoeba's complement is $\text{amoeba}(R)^c = \mathbb{R}^d \setminus \text{amoeba}(R)$.

The following result follows from Gelfand, Kapranov, and Zelevinsky [6, Chap. 6, Prop. 1.5 and Cor. 1.6].

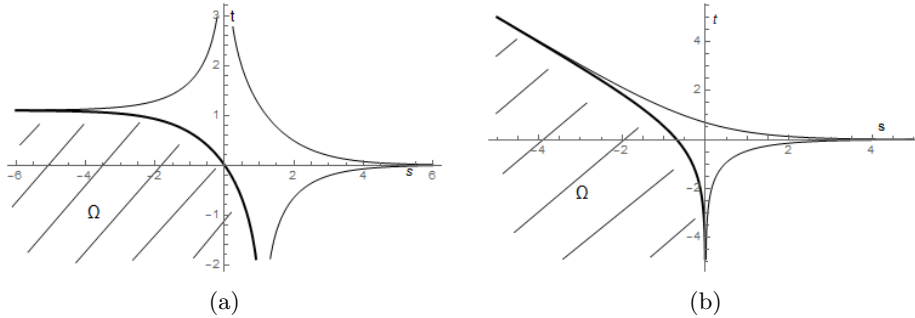


Figure 6: The $\text{amoeba}(f_1)$, $\text{amoeba}(f_2)$ along with sets Ω associated to the language L_{F_2} and the Fibonacci language are presented in 6a and 6b, respectively, where $f_1(\mathbf{z}) = 1 - z_1 - z_2 - 3z_1z_2$ and $f_2(\mathbf{z}) = 1 - z_1 - z_1z_2$.

Proposition 5.1. *If $f(\mathbf{z})$ is a Laurent polynomial then all connected components of the complement $\text{amoeba}(f)^c$ are convex subsets of \mathbb{R}^d . These real convex subsets are in bijection with the Laurent series expansions of the rational function $\frac{1}{f(\mathbf{z})}$. When $\frac{1}{f(\mathbf{z})}$ has a power series expansion, then it corresponds to the component of $\mathbb{R}^d \setminus \text{amoeba}(f)$ containing all points $(-N, \dots, -N)$ for N positive and sufficiently large.*

The techniques of drawing amoeba are well developed. The amoebas of polynomials $f_1(\mathbf{z}) = 1 - z_1 - z_2 - 3z_1z_2$ and $f_2(\mathbf{z}) = 1 - z_1 - z_1z_2$ that correspond to the cases of language L_{F_2} of freely reduced words of F_2 and Fibonacci language, respectively, are presented in the Figure 6.

Looking at the shape of $-\Omega$ presented in Figure 7, in the case of L_{F_2} , we see that for each $\mathbf{r} \in M_2$ with positive entries there is a tangent line to the boundary $\partial(-\Omega)$, is orthogonal to \mathbf{r} and hence, the infimum of the linear form (23) is achieved. To compute $\psi(\mathbf{r})$ we apply a standard method of Lagrange multipliers.

In coordinates $(s, t) \in \mathbb{R}^2$, $z_1 = e^{-s}$, $z_2 = e^{-t}$, the boundary of $-\Omega$ is a curve l given by the equation

$$1 - e^{-s} - e^{-t} - 3e^{-s-t} = 0 \quad (28)$$

or

$$e^{s+t} - e^s - e^t - 3 = 0 \quad (29)$$

Let $\mathbf{r} = (p, q) \in M_2$. We have to minimize $ps + qt$ when $(s, t) \in l$. The associated Lagrange function is

$$\Phi(s, t, \lambda) = ps + qt - \lambda(e^{s+t} - e^s - e^t - 3).$$

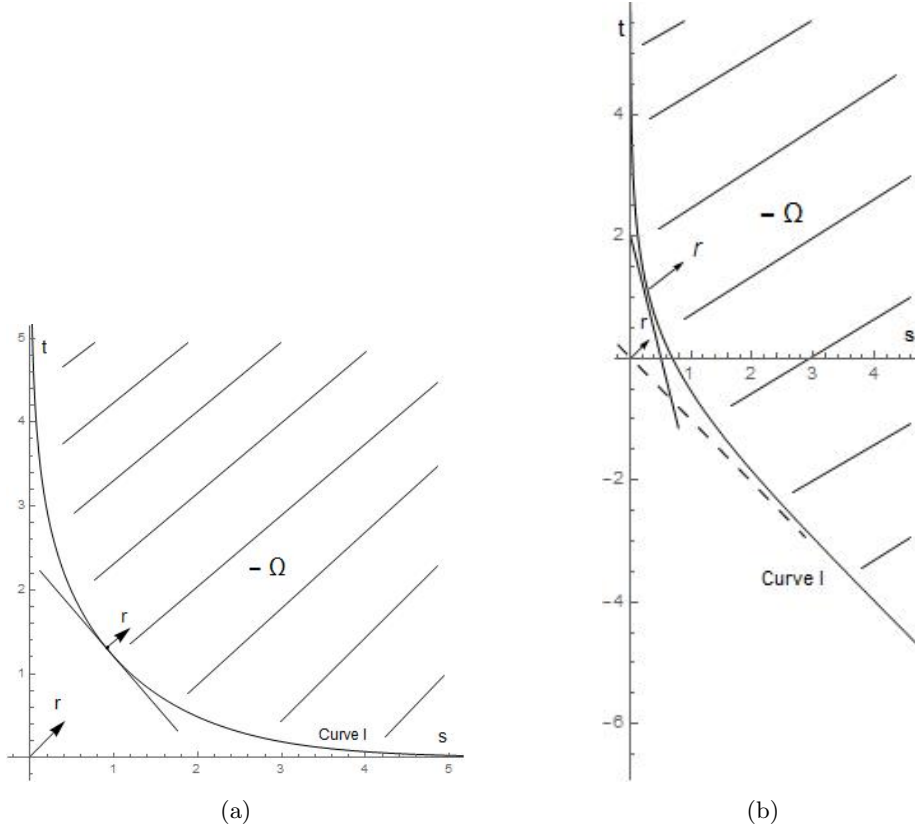


Figure 7: The sets $-\Omega$ presented in 7a and 7b are associated with languages L_{F_2} and L_{Fib} , respectively.

Equating partial derivatives to zero we obtain

$$\begin{aligned} \frac{\partial \Phi}{\partial s} &= p - \lambda(e^{s+t} - e^s) = 0 \implies p = \lambda(e^{s+t} - e^s) \\ \frac{\partial \Phi}{\partial t} &= q - \lambda(e^{s+t} - e^t) = 0 \implies q = \lambda(e^{s+t} - e^t) \\ \frac{\partial \Phi}{\partial \lambda} &= e^{s+t} - e^s - e^t - 3 = 0. \end{aligned}$$

This gives

$$\rho = \frac{p}{q} = \frac{1 - e^{-t}}{1 - e^{-s}} \implies e^{-t} = 1 - (1 - e^{-s})\rho$$

Substituting the value of e^{-t} in (28) we get

$$\begin{aligned} 1 - e^{-s} - (1 - (1 - e^{-s})\rho) - 3e^{-s} (1 - (1 - e^{-s})\rho) &= 0 \\ -e^{-s} + \rho - \rho e^{-s} - 3e^{-s} + 3\rho e^{-s} - 3\rho e^{-2s} &= 0 \\ -3\rho e^{-2s} + (2\rho - 4)e^{-s} + \rho &= 0 \end{aligned}$$

Substituting $x = e^{-s}$ in the above quadratic and then solving it gives

$$x_{1,2} = \frac{2\rho \pm \sqrt{1 - \rho + \rho^2}}{6\rho}$$

We choose positive sign (i.e. +) because we know due to [9] that the function is real analytic. Re-substituting $x = e^{-s}$ and $\rho = \frac{p}{q}$ we get

$$e^{-s} = \frac{p - 2q + 2\sqrt{p^2 - pq + q^2}}{3p},$$

$$e^s = \frac{2q - p + 2\sqrt{p^2 - pq + q^2}}{p}$$

and by symmetry,

$$e^t = \frac{2p - q + 2\sqrt{p^2 - pq + q^2}}{q}.$$

Hence,

$$s = \log \left(\frac{2q - p + 2\sqrt{p^2 - pq + q^2}}{p} \right), \quad t = \log \left(\frac{2p - q + 2\sqrt{p^2 - pq + q^2}}{q} \right),$$

and

$$\psi_{F_2}(\mathbf{r}) = p \log \left(\frac{2q - p + 2\sqrt{p^2 - pq + q^2}}{p} \right) + q \log \left(\frac{2p - q + 2\sqrt{p^2 - pq + q^2}}{q} \right)$$

or

$$\psi_{F_2}(\mathbf{r}) = \mathbf{H}(\mathbf{r}) + p \log \left(2q - p + 2\sqrt{p^2 - pq + q^2} \right) + q \log \left(2p - q + 2\sqrt{p^2 - pq + q^2} \right),$$

where $\mathbf{H}(\mathbf{r}) = -p \log p - q \log q$. See Figure 2a for the graph of $\psi_{F_2}(\mathbf{r}) = \psi_{F_2}(p, 1 - p)$.

6 Multivariate growth in the case of Fibonacci language

Fibonacci language L_{Fib} is a language over binary alphabet $\{0, 1\}$ consisting of words that do not contain 11 as a subword i.e. the word 11 is forbidden. It is

one of the important examples associated with subshifts of finite type [15]. The reason why Fibonacci name is associated to it is coming from the fact that the number of words of length n in L_{Fib} is equal to the $(n+2)$ th Fibonacci number β .

$$\beta = \frac{1}{\sqrt{5}} (\lambda_1^{n+2} - \lambda_2^{n+2}),$$

where $\lambda_1 = \frac{1 + \sqrt{5}}{2}$ and $\lambda_2 = \frac{1 - \sqrt{5}}{2}$ are eigenvalues of matrix

$$A = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$$

which is the matrix of transitions of the Fibonacci subshift. At the same time A is the adjacency matrix of the automaton \mathcal{A}_{Fib} given in the Figure 8 which

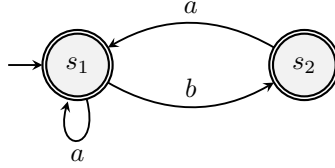


Figure 8: Moore diagram of Fibonacci automaton \mathcal{A}_{Fib}

accepts the language L_{Fib} , if instead of $\{0, 1\}$ we use the alphabet $\{a, b\}$. An easy computation of growth series gives function

$$\Gamma_{Fib}(\mathbf{z}) = \frac{1 + z_2}{1 - z_1 - z_1 z_2}$$

mentioned before Definition 2.1. In order to compute ψ , observe from Figure 8 that the Condition (CG) holds as the automaton \mathcal{A}_{Fib} is ergodic.

Proposition 6.1. *The indicatrice ψ_{Fib} viewed as a function of $p, 0 < p < 1$ where $\mathbf{r} = (p, 1 - p) \in M_2$ is the direction vector, is given by*

$$\psi_{Fib}(\mathbf{r}) = \begin{cases} p \log \left(\frac{p}{2p-1} \right) + (1-p) \log \left(\frac{2p-1}{1-p} \right) & \text{if } p \geq \frac{1}{2} \\ -\infty & \text{if } p < \frac{1}{2} \end{cases}$$

The graph of $\psi_{Fib}(\mathbf{r})$ on $[\frac{1}{2}, \infty)$ is shown in the Figure 2b.

Proof. As before, we switch to the variables x, y instead of z_1, z_2 , respectively. The amoeba of $f_2(x, y) = 1 - x - xy$ and the set $-\Omega$ are shown in the Figures 6b and 7b, respectively.

We provide two explanations why $\psi_{Fib}(p) = -\infty$ when $0 < p < \frac{1}{2}$, first, algebraic and second, geometric. The power series expansion of $(1 - x - xy)^{-1}$ is

$$\begin{aligned} \frac{1}{1 - x - xy} &= \sum_{n=1}^{\infty} \sum_{i=0}^n \binom{n}{i} x^{n-i} (xy)^i \\ &= \sum_{n=1}^{\infty} x^n \sum_{i=0}^n \binom{n}{i} y^i \end{aligned}$$

Hence,

$$\begin{aligned} \frac{1+y}{1-x-xy} &= \sum_{n=1}^{\infty} x^n \sum_{i=0}^n \binom{n}{i} y^i + \sum_{n=1}^{\infty} x^n \sum_{i=0}^n \binom{n}{i} y^{i+1} \\ &= \sum_{n=1}^{\infty} x^n \left\{ 1 + y + \sum_{i=0}^n \left[\binom{n}{i} + \binom{n}{i-1} \right] y^i \right\} \end{aligned}$$

and we see that coefficients of the power series corresponding to the indices (n, i) with $i > n + 1$ are zero. Hence, any direction $\mathbf{r} = (p, q)$ with $\frac{p}{q} > 1$ gives value $-\infty$ to ψ_{Fib} . As for the geometric explanation let us use Figures 6b and 7b. The domain $-\Omega$ is shown in Figure 7b and it is bounded by the curve l given by the equation

$$1 - e^{-s} - e^{-s-t} = 0 \quad (30)$$

- Case (i) If $p < \frac{1}{2}$, then $\frac{1-p}{p} > \frac{1}{2}$ and lines

$$ps + (1-p)t = c$$

intersect $-\Omega$ for arbitrary $c \in \mathbb{R}$. Hence,

$$\psi_{Fib}(p) = \inf_{(s,t) \in -\Omega} (ps + (1-p)t) = -\infty$$

- Case (ii) If $p > \frac{1}{2}$, then $\frac{1-p}{p} < \frac{1}{2}$ and there is a unique value of c such that the line

$$ps + (1-p)t = c$$

is tangent to the curve l . The coordinates (s_0, t_0) of the tangent point P gives a minimum value to the linear form $ps + (1-p)t$ when $(s, t) \in -\Omega$. To find it again, we again apply the method of Lagrange. We denote $q = 1-p$ and rewrite Equation (30) as

$$e^{s+t} - e^t - 1 = 0$$

The associated Lagrange function is

$$\Phi(s, t, \lambda) = ps + qt - \lambda(e^{s+t} - e^t - 1)$$

and the corresponding system of equation is

$$\begin{aligned}\frac{\partial\Phi}{\partial s} &= p - \lambda e^{s+t} = 0 \\ \frac{\partial\Phi}{\partial t} &= q - \lambda e^{s+t} + \lambda e^t = 0 \\ \frac{\partial\Phi}{\partial\lambda} &= e^{s+t} - e^t - 1 = 0\end{aligned}\tag{31}$$

From the first equation in (31), we get $e^{s+t} = \frac{p}{\lambda}$. from the second equation we get $q = p + \lambda e^t$, and from the third equation we get $\frac{p}{\lambda} - e^t - 1 = 0$. Hence,

$$e^t = \frac{p - q}{\lambda} = \frac{2p - 1}{\lambda}$$

But $\lambda = 1 - p = q$. This gives,

$$e^t = \frac{2p - 1}{1 - p} \text{ and hence, } e^s = \frac{p}{2p - 1}.$$

The last equation determine the point (s_0, t_0) . Making substitution in (23) we obtain (30).

□

7 Multivariate growth and LDT

The alternative approach for computing $\psi(\mathbf{r})$ is via the application of methods of Large Deviation Theory (LDT). Here we discuss a special case related to languages associated with subshifts of finite type. Let us first recall the basic facts about subshifts of finite type (SFT). For more details see [15]. Let $\Sigma = \{a_1, \dots, a_d\}$ be a finite alphabet and $\Sigma^{\mathbb{Z}}$ be a space of two-sided infinite sequences over Σ indexed by integers. $\Sigma^{\mathbb{Z}}$ is supplied by a product topology and is homeomorphic to a Cantor set when $d \geq 2$. The shift map $U : \Sigma^{\mathbb{Z}} \rightarrow \Sigma^{\mathbb{Z}}$ is the homeomorphism given by

$$(Uw)_n = w_{n+1}, \quad w = (w)_{n=-\infty}^{\infty} \in \Sigma^{\mathbb{Z}}.$$

A closed U -invariant subset $X \subset \Sigma^{\mathbb{Z}}$ is a subshift. Let $L_X \subset \Sigma^*$ be a language of subshift consisting of (finite) words that appear as a subwords of $w \in X$. A subshift X is said to be subshift of finite type if there is a finite subset $\mathbf{F} \subset \Sigma^*$ (set of forbidden words) such that $X = X_{\mathbf{F}}$ consist of sequences $w \in \Sigma^{\mathbb{Z}}$ that do not contain forbidden subwords. It is obvious that $X_{\mathbf{F}}$ is closed and U -invariant. For instance, in the case $\Sigma = \{0, 1\}$ and $\mathbf{F} = \{11\}$ we get the Fibonacci subshift. Alternative way to define subshifts of finite type is via the graph $\Xi = (V, E)$ (in fact directed multi-graph i.e. loops and multiple edges are allowed) or equivalently, via the adjacency matrices $A = (a_{ij})$ of the size $|V| \times |V|$ whose rows and columns correspond to vertices of the graph and $a_{ij}, i, j \in V$ is

a non-negative integer equal to the number of edges joining vertex i with vertex j . For instance, the Fibonacci subshift the matrix is

$$A = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$$

and the graph Ξ is given by the Figure 9 and is similar to the diagram of the automaton from Figure 8.

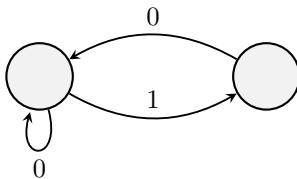


Figure 9: The graph Ξ of the Fibonacci subshift

Another example is the free group F_m , the subshift X , $m \geq 2$ with alphabet $\Sigma = \{a_1, \dots, a_m, a_1^{-1}, \dots, a_m^{-1}\}$ and

$$A = \begin{pmatrix} J_{m \times m} & K_{m \times m} \\ K_{m \times m} & J_{m \times m} \end{pmatrix},$$

where J is matrix with all entries 1, $K = J - I$ and I is an $m \times m$ identity matrix. The language L_X consists of freely reduced words.

It is well known that the language L_X associated with a subshift of finite type is regular. Hence, its multivariate growth series represents a rational function and the technique of computation of multivariate growth rate described in the previous sections is applicable. Now, we shall see how the results of LDT can be used for the same goal.

To make one more step towards LDT, let us recall some other important notions related to SFT. The SFT (U, X) is *irreducible* if the graph Ξ_X is strongly connected (i.e. for any vertices in graph, there is a path connecting them). In this case, the associated matrix A is called irreducible. The Perron-Frobenius theorem states that the irreducible matrix A with non-negative entries (like in our case) has a simple eigenvalue $\rho = \rho(A)$ (called Perron-Frobenius eigenvalue) such that any other eigenvalue λ satisfies $|\lambda| \leq \rho$. Also there are two vectors $\mathbf{u}, \mathbf{v} \in \mathbb{R}^d$ satisfying

$$A\mathbf{v} = \rho\mathbf{v}, \mathbf{u}^t A = \rho\mathbf{u}^t,$$

where \mathbf{u}^t is the transpose of the column vector \mathbf{u} . These vectors are unique up to scalar factor. Assume that A is primitive matrix. Then the Perron-Frobenius triple $(\rho, \mathbf{u}, \mathbf{v})$ (consisting of vectors $\mathbf{u}, \mathbf{v} > 0$ such that $A\mathbf{v} = \rho\mathbf{v}, \mathbf{u}^t A = \rho\mathbf{u}^t$, and normalized by the condition $\langle \mathbf{u}^t, \mathbf{v} \rangle = 1$) gives the information about the

powers A^n of A :

$$\lim_{n \rightarrow \infty} \frac{1}{\rho^n} A^n = \mathbf{v} \cdot \mathbf{u}^t.$$

See ([15], Theorem 4.5.12).

Next, we need the notion of Markov measure on $\Sigma^{\mathbb{Z}}$. Given a stochastic $d \times d$ matrix

$$P = (p_{ij}), \text{ where } p_{ij} \geq 0 \text{ and } \sum_{j=1}^d p_{ij} = 1, i = 1, \dots, d,$$

and a stationary probability row vector $p = (p_1, \dots, p_d)$, $pP = p$, one can define a Borel probability measure $\mu = \mu_P$ on $\Sigma^{\mathbb{Z}}$ by

$$\mu([\omega_0, \dots, \omega_n]) = p_{\omega_0} \prod_{i=0}^{n-1} p_{\omega_i, \omega_{i+1}},$$

where $[\omega_0, \dots, \omega_n]$ is a cylinder subset of $\Sigma^{\mathbb{Z}}$ consisting of all $\omega \in \Sigma^{\mathbb{Z}}$ with the prescribed entries $\omega_0, \dots, \omega_n$ at coordinates $0, 1, \dots, n$. (such sets generate the sigma-algebra of Borel subsets and hence, values of μ on them determine μ completely). The Perron-Frobenius eigenvalue of P is 1, vector p exists and is unique if P is irreducible. The measure μ_P is shift-invariant and the system $(\Sigma^{\mathbb{Z}}, U, \mu)$ is ergodic.

One of the main results of theory of SFT is a statement (assuming irreducibility of X) on the existence and uniqueness of probability measure $\eta = \eta_X$ of maximal entropy. Not getting into the details we just mentioned that if ϑ is U -invariant probability measures on $\Sigma^{\mathbb{Z}}$, then the Kolmogorov-Sinai entropy $h(\vartheta)$ can be defined. Then

$$h(\eta) = \max_{\vartheta} h(\vartheta),$$

where maximum is taken over all U -invariant probability measures supported on X . The measure η is called *Parry* measure. It is a Markov type measure determined by a stochastic matrix $P = (p_{ij})$, with

$$p_{ij} = \frac{1}{\rho} a_{ij} \frac{v_j}{v_i}, \tag{32}$$

where $A = (a_{ij})$ and $A\mathbf{v} = \rho\mathbf{v}$. For such measure we have

$$\mu([i, x_1, \dots, x_{n-1}, j]) = \frac{u_i v_j}{\rho^n} \tag{33}$$

See [17] for further details. In fact, for (32) and (33) we have to assume that $a_{ij} \in \{0, 1\}$ (i.e. graph Γ does not have multiple edges). Meanwhile observe that, every SFT can be coded in such a way that the matrix A will have only

entries $\{0, 1\}$ [15]. Starting from this moment we assume that the measure of maximal entropy is associated with SFT.

Now, we recall some basic notions of LDT, namely, the notion of the rate function I and the Large Deviation Principal (LDP). A rate function is a lower semi-continuous function $I : W \rightarrow [0, +\infty]$ defined on a topological space W (for us now $W = \mathbb{R}^d$) such that for each $a \in [0, \infty)$ the level set

$$Y_I(a) = \{\mathbf{x} \in W : I(\mathbf{x}) \leq a\}$$

is a closed subset of W . A *good rate function* is a rate function for which all level sets $Y_I(a)$ are compact. A sequence $\{\mu_n\}_{n=1}^\infty$ of Borel measures on W satisfies LDP if, for every Borel subset $B \subset W$,

$$-\inf_{\mathbf{x} \in B^\circ} I(\mathbf{x}) \leq \liminf_{n \rightarrow \infty} \frac{1}{n} \log \mu_n(B) \leq \limsup_{n \rightarrow \infty} \frac{1}{n} \log \mu_n(B) \leq -\inf_{\mathbf{x} \in \overline{B}} I(\mathbf{x}),$$

where B°, \overline{B} are, respectively, the interior and the closure set of the set B .

Given a finite Markov chain on a set $\Sigma = \{a_1, \dots, a_d\}$ determined by a stochastic matrix $P = (p_{ij})_{i,j}^d$ and a function $f : \Sigma \rightarrow \mathbb{R}^d, d \geq 1$ one can consider for each $x = (x_i)_{i=1}^\infty \in \Sigma^\mathbb{N}$ the empirical means

$$Z_n(x) = \frac{1}{n} \sum_{i=1}^n f(x_i)$$

and the corresponding distributions μ_n (discrete measures in \mathbb{R}^d). We assume that the random process is a Markov process given by the matrix P and stationary vector $p, pP = p$.

Associate with every $\mathbf{y} \in \mathbb{R}^d$ a non-negative matrix $\Pi(\mathbf{y})$ whose elements are

$$\pi_{ij}(\mathbf{y}) = p_{ij} e^{\langle \mathbf{y}, f(j) \rangle}, i, j \in \Sigma.$$

$\Pi(\mathbf{y})$ is a matrix with non-negative entries and it is irreducible if and only if P is irreducible. Let $\rho(\Pi(\mathbf{y}))$ be the Perron Frobenius eigenvalue of $\Pi(\mathbf{y})$. The Theorem 3.1.2 of [5] states that the empirical mean Z_n (or corresponding distributions μ_n) satisfies the LDP with the convex and good rate function

$$I(\mathbf{z}) = \sup_{\mathbf{y} \in \mathbb{R}^d} \{\langle \mathbf{y}, \mathbf{z} \rangle - \log \rho(\Pi(\mathbf{y}))\}, \mathbf{z} \in \mathbb{R}^d.$$

There is a version of this result due to Sanov which is more suitable for our goals. Let $f : \Sigma \rightarrow \mathbb{R}^d, |\Sigma| = d$ be a function such that

$$f(a_i) = \begin{pmatrix} 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{pmatrix}$$

where 1 is at i -th position and $i = 1, \dots, d$.

The following alternative description of I holds ([5], Theorem 3.1.6).

$$I(\mathbf{r}) = \begin{cases} \sup_{\mathbf{u} \gg \mathbf{0}} \sum_{j=1}^d r_j \log \left(\frac{u_j}{(\mathbf{u}P)_j} \right) & \mathbf{r} \in M_d \\ \infty & \mathbf{r} \notin M_d \end{cases}$$

The supremum is taken over the strictly positive vectors \mathbf{u} i.e. $u_i > 0$ for all i .

Now we have all needed to describe the connection between ψ and I . We assume that the language $L \subset \Sigma^*$ is a language determined by the automaton \mathcal{A} with the property that for each state $q \in Q$ all incoming edges are labeled by the same symbol (like in the examples given by the Figure 3 or 8). The indicatrice of growth as before is denoted by $\psi(\mathbf{r})$. We assume that \mathcal{A} is ergodic (and hence, condition (CG) satisfied). Let A be the adjacency matrix of \mathcal{A} and $X_{\mathcal{A}}$ be the corresponding subshift. Let L_1 be a language associated with $X_{\mathcal{A}}$. In described situation we identify Σ with Q attaching to each state $q \in Q$ symbol a_q (the label of the entering edges). Because \mathcal{A} is ergodic, we get a bijection between Σ and Q . After such identification, it become obvious that $L \subset L_1$. Because of ergodicity (i.e. irreducibility of A) as easily can be shown the indicatrice of growth $\psi(\mathbf{r})$ is same for languages L and L_1 .

Associate with $A = (a_{ij})$ the stochastic matrix $P = (p_{ij})$, where $p_{ij} = \frac{a_{ij} v_j}{\rho v_i}$, $\rho = \rho(A)$ and $A\mathbf{v} = \rho\mathbf{v}$, $\mathbf{v} \gg \mathbf{0}$. Let p be a stationary probability vector ($pP = p$) and $\mu = \mu_p$ be a corresponding Markov measure (i.e. Perry measure). From (34) (assuming normalization $\langle p, \mathbf{v} \rangle = 1$) we know that the measure μ is almost equidistributed on the cylinder sets C_w determined by the words $w \in L_1$ of the fixed length as for any $i, j, w_1, \dots, w_{n-1} \in \Sigma$

$$\mu([i, w_1, \dots, w_{n-1}, j]) = \frac{p_i v_j}{\rho^n}. \quad (34)$$

Let $\mathbf{r} \in M_d$ be a rational vector with positive entries and $\mathcal{C} \subset M_d$ a small neighborhood of \mathbf{r} . Let

$$B_n = \{w \in \Delta_{\mathcal{A}} : Z_n(w) \in \mathcal{C}\}$$

From LDP, we know that $\frac{1}{n} \log \mu(B_n)$ is close to $-I(\mathbf{r})$ when n is large. On the other hand, from (34) we get that $\frac{1}{n} \log \mu(B_n)$ is close to $\frac{1}{n} \log(\rho^{-n} \cdot l_{n\mathbf{r}})$, where

$$L_1(\mathbf{z}) = \sum_{\mathbf{i} \in \mathbb{N}^d} l_{\mathbf{i}} \mathbf{z}^{\mathbf{i}}$$

is a multivariate growth series of L_1 . Hence, in the limit when $n \rightarrow \infty$ we get

the equality

$$\begin{aligned} I(\mathbf{r}) &= \log \rho - \limsup_{n \rightarrow \infty} \frac{1}{n} \log l_{n\mathbf{r}} \\ &= \log \rho - \psi(\mathbf{r}) \end{aligned} \quad (35)$$

In fact, under the impose conditions using the definition of $\psi(\mathbf{r})$ one can prove (35) for all $\mathbf{r} \in M_d$, not just rationals. We summarize this as

Proposition 7.1. *The indicatrice of growth $\psi(\mathbf{r})$ of a language determined by ergodic automaton \mathcal{A} of type described above satisfies*

$$I(\mathbf{r}) = \log \rho - \psi(\mathbf{r}),$$

where $I(\mathbf{r})$ is the rate function associated with the Markov chain determined by the stochastic matrix P corresponding to A and

$$I(\mathbf{r}) = \sup_{\mathbf{u} \gg \mathbf{0}} \sum_{j=1}^d r_j \log \left[\frac{u_j}{(\mathbf{u}P)_j} \right], \quad \mathbf{r} \in M_d. \quad (36)$$

Finally, we make one more remark. Recall that entries of A belong to the set $\{0, 1\}$. Let as before $\mathbf{v} = (v_1, \dots, v_d)$ be a right eigenvector of A corresponding to $\rho = \rho(A)$, $A\mathbf{v} = \rho\mathbf{v}$. Assume that \mathbf{v} is a probability vector. Let

$$M_{d>0} = \{\mathbf{r} = (r_1, \dots, r_d) \in M_d : r_i > 0, \forall i\}.$$

Then the map $T : M_{d>0} \rightarrow \mathbb{R}^d$ given by

$$T(\mathbf{q}) = \mathbf{s} = (s_1, \dots, s_d),$$

where

$$s_j = \frac{q_j}{v_j \sum_{i=1}^d \frac{q_i a_{ij}}{v_i}}, \quad j = 1, \dots, d. \quad (37)$$

Let $A(\mathbf{z})$ be the matrix from Proposition 2.1 i.e. matrix obtained from A by replacing each 1 in the j th column of A by z_j and let \mathbf{t} be a vector obtained from $\mathbf{q} \in M_{d>0}$ by

$$\mathbf{t} = \left(\frac{q_1}{v_1}, \dots, \frac{q_d}{v_d} \right)$$

Lemma 7.1. *For each $\mathbf{q} \in M_{d>0}$ vector \mathbf{t} satisfies $\mathbf{t}A(\mathbf{s}) = \mathbf{t}$, where $\mathbf{s} = (s_1, \dots, s_d)$ is given by (37).*

Proof. For $1 \leq j \leq d$,

$$\begin{aligned} (\mathbf{t}A(\mathbf{s}))_j &= s_j \sum_{i=1}^d \frac{q_i a_{ij}}{v_i} \\ &= \frac{q_j}{v_j \sum_{i=1}^d \frac{q_i a_{ij}}{v_i}} \cdot \sum_{i=1}^d \frac{q_i a_{ij}}{v_i} \\ &= \frac{q_j}{v_j} = t_j \end{aligned} \quad (38)$$

(in the above relations we used the fact that $a_{ij} \in \{0, 1\}$). \square

Observe that \mathbf{s} depends on the vector \mathbf{q} , so we can write $\mathbf{s} = \mathbf{s}(\mathbf{q})$.

Corollary 7.1.

$$\det(I - A(\mathbf{s}(\mathbf{q}))) = 0, \quad \forall \mathbf{q} \in M_{d>0}.$$

Recall that the multivariate growth series of language L satisfies Proposition 2.1. i.e.

$$\Gamma_L(\mathbf{z}) = \frac{G(\mathbf{z})}{\det(I - A(\mathbf{z}))}$$

and singularities of $\Gamma_L(\mathbf{z})$ are determined by the roots of denominator. Hence, $T(M_{d>0})$ is a part of the set of the real singularities of $\Gamma_L(\mathbf{z})$ and is part of the boundary $\partial\mathcal{D}$.

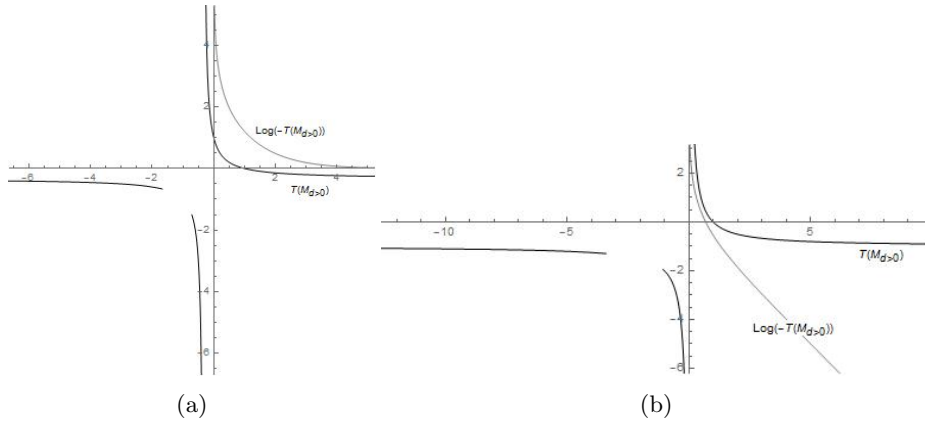


Figure 10: The graphics of the sets $T(M_{d>0})$ and $\log(-T(M_{d>0}))$ presented in (10a) and (10b) are associated with languages L_{F_2} and L_{Fib} , respectively.

We already know from Proposition 7.1 that

$$I(\mathbf{r}) = \log \rho - \psi(\mathbf{r}).$$

Recall that $P = (p_{ij})$ and $p_{ij} = \frac{v_j a_{ij}}{\rho v_i}$. Rewriting (36) as

$$\begin{aligned} I(\mathbf{r}) &= \sup_{\mathbf{q} \in M_{d>0}} \left\{ \sum_{j=1}^d r_j \log \rho - \sum_{j=1}^d r_j \log \left(\frac{v_j \sum_{i=1}^d \frac{q_i a_{ij}}{v_i}}{q_j} \right) \right\} \\ &= \log \rho - \inf_{\mathbf{q} \in M_{d>0}} \sum_{j=1}^d r_j (-1) \log [T(\mathbf{q})]_j \\ &= \log \rho - \inf_{\mathbf{s} \in T(M_d)} \left(- \sum_{j=1}^d r_j \log s_j \right) \end{aligned}$$

we conclude

$$\psi(\mathbf{r}) = \inf_{\mathbf{s} \in T(M_d)} \left(- \sum_{j=1}^d r_j \log s_j \right)$$

Comparing with (26) we observe that in the described situation, the infimum in (26) can be taken only via the subset of the positive part of the boundary $\partial\mathcal{D}$. See Figures 10.

8 Finer asymptotic

As was already mentioned in the introduction, in the case of rational vector \mathbf{r} the function $\psi(\mathbf{r})$ can be identified by (3) (alternatively, the function $\varphi(\mathbf{r}) = e^{\psi(\mathbf{r})}$ can be identified by (14)). A powerful results of the theory of ACSV (asymptotic combinatorics in several variables) presented in [16] allow not only to compute in many cases $\varphi(\mathbf{r})$, for rational functions $\Gamma(\mathbf{z}) = \frac{G(\mathbf{z})}{H(\mathbf{z})}$ and rational vectors \mathbf{r} but also to describe much finer asymptotics of the diagonal coefficients $f_{n\mathbf{r}}$ (See Theorems 5.1,5.2,5.3 in [16]). Also there are statements on the smoothness of $\varphi(\mathbf{r})$ as a function of rational \mathbf{r} .

Our definition of $\varphi(\mathbf{r})$, that follow the idea from [18] and is based on the cone approach works for arbitrary direction $\mathbf{r} \in M_d$. Moreover, under the assumption (CG) in many “good” examples (including considered in this paper) the upper bound (17) can be replaced by the equality

$$\varphi(\mathbf{r}) = e^{\inf_{\mathbf{x} \in \overline{\mathcal{D}}} h_{\mathbf{r}}(\mathbf{x})} = e^{\inf_{\mathbf{x} \in \partial\mathcal{D}} h_{\mathbf{r}}(\mathbf{x})} = e^{\psi(\mathbf{r})}$$

when \mathbf{r} is rational, where

$$\psi(\mathbf{r}) = \inf_{\theta \in \partial(-\overline{\Omega})} \langle \mathbf{r}, \theta \rangle, \quad \Omega = \text{Relog}(\mathcal{D})$$

Additionally, the facts based on the convex analysis and Large Deviation Theory allow to claim that for the good rational functions $\Gamma(\mathbf{z})$, the function $\varphi(\mathbf{r})$ is a real analytic function. See [9].

Now we recall few definitions and results presented in [16], apply them to our examples and make a comparison. The Theorem 5.1 from [16] basically states the following.

Let $\mathbf{r} \in \mathbb{Q}^d$ and let $G(\mathbf{z}), H(\mathbf{z}) \in \mathbb{Q}[\mathbf{z}]$ be coprime polynomials such that

$$\Gamma(\mathbf{z}) = \frac{G(\mathbf{z})}{H(\mathbf{z})}$$

admits a power series expansion

$$\Gamma(\mathbf{z}) = \sum_{\mathbf{i} \in \mathbb{N}^d} f_{\mathbf{i}} \mathbf{z}^{\mathbf{i}}.$$

Suppose that the system of equations

$$H(\mathbf{z}) = r_2 z_1 H_{z_1}(\mathbf{z}) - r_1 z_2 H_{z_2}(\mathbf{z}) = r_d z_1 H_{z_1}(\mathbf{z}) - r_1 z_d H_{z_d}(\mathbf{z}) \quad (39)$$

admits a finite number of solutions, exactly one of which $\mathbf{z}' \in \mathbb{C}_*^d$ is minimal (i.e. no other singularity \mathbf{z} of $\Gamma(\mathbf{z})$ satisfies $|z_j| < |z'_j|$ for all $1 \leq j \leq d$).

Suppose that $H_{z_d}(\mathbf{z}) \neq 0, G(\mathbf{z}') \neq 0$. Then as $n \rightarrow \infty$

$$f_{n\mathbf{r}} = \mathbf{z}'^{-n\mathbf{r}} n^{\frac{1-d}{2}} \frac{(2\pi r_d)^{\frac{1-d}{2}} - G(\mathbf{z}')}{z'_d H_{z_d}(\mathbf{z}') \sqrt{\det(\mathcal{H})}} \left(1 + O\left(\frac{1}{n}\right) \right) \quad (40)$$

when $n\mathbf{r} \in \mathbb{N}^d$, where \mathcal{H} is a $(d-1) \times (d-1)$ matrix defined by Equation (5.25) in [16] and it is supposed that $\det(\mathcal{H}) \neq 0$.

It is also claimed that as \mathbf{r} varies in any sufficiently neighborhood in $\mathbb{R}_{>0}^d$, the solution $\mathbf{z}' = \mathbf{z}'(\mathbf{r})$ varies smoothly with \mathbf{r} . The factor $(\mathbf{z}')^{-n\mathbf{r}}$ in (40) gives us

$$\limsup_{n \rightarrow \infty} |f_{n\mathbf{r}}|^{\frac{1}{n}} = |(\mathbf{z}')^{-\mathbf{r}}| = e^{-\sum r_i \log |z'_i|} = e^{-h_{\mathbf{r}}(\mathbf{z}')}$$

The system of equations (39) is equivalent in our situation (assuming condition (CG)) to the system of equations coming from the Lagrange multipliers method because if

$$\Phi(\mathbf{z}, \lambda) = \sum_{i=1}^d r_i \log z_i - \lambda H(\mathbf{z}) \quad (41)$$

then the critical points are solutions of the system

$$\begin{aligned} \frac{\partial \Phi}{\partial z_i} &= \frac{r_i}{z_i} - \lambda H_{z_i} = 0, i = 1, \dots, d \\ H(\mathbf{z}) &= 0 \end{aligned} \quad (42)$$

which is equivalent to (39). In our situation, we make substitutions $z_i = e^{-\theta_i}$ and replace (41) by the

$$\Phi'(\theta, \lambda) = \sum_{i=1}^d r_i \theta_i - \lambda H(e^{-\theta})$$

as shown in the previous sections. Of course (40) is much finer asymptotic than

$$f_{n\mathbf{r}} \sim (\mathbf{z}')^{-n\mathbf{r}}.$$

As was already mentioned, the condition (CG) gives an alternative definition of the growth in the direction of \mathbf{r} that works not only for rational \mathbf{r} but for

any direction $\mathbf{r} \in \mathbb{R}_{>0}^d$. At the same time the above argument show that when (CG) holds and conditions of Theorem 5.1 in [16] are satisfied, then the rates of growth in the rational direction \mathbf{r} defined by (3) or as $e^{-\psi(\mathbf{r})}$ coincide. The smooth dependence of $\mathbf{z}' = \mathbf{z}'(\mathbf{r})$ (and hence of $\varphi(\mathbf{r}) = \mathbf{z}'^{-\mathbf{r}}$) on $\mathbf{r} \in \mathbb{Q}_{>0}^d$ can be strengthened to the claim about the analytic dependence on \mathbf{r} for all $\mathbf{r} \in \mathbb{R}_{>0}^d$, where $\psi(\mathbf{r}) \geq 0$ as shown in [9].

An useful tool in discussed topics is the logarithmic gradient map

$$\nabla_{\log} f = (z_1 f_{z_1}, \dots, z_d f_{z_d}).$$

Proposition 3.13 in [16] states that for any minimal singular point \mathbf{z}' of $\Gamma(\mathbf{z}) = \frac{G(\mathbf{z})}{H(\mathbf{z})}$ (where G, H are coprime) there exists $\mathbf{r} \in \mathbb{R}_{\geq 0}^d$ and $\tau \in \mathbb{C}$ such that

$$\left(\nabla_{\log} H^{(\mathbf{s})} \right) (\mathbf{z}') = \tau \cdot \mathbf{r},$$

where $H^{(\mathbf{s})}$ is a square free part of H . In this situation \mathbf{z}' is either a minimizer or a maximizer of the height function $h_{\mathbf{r}}(\mathbf{z})$ on \mathcal{D} .

Let us apply Theorem 5.1 from [16] to the case of the free group of rank 2. We assume that $\mathbf{r} = (p, 1-p)$ is rational, $\mathbf{z} = (x, y)$. Recall that the multivariate growth series of free group F_2 is

$$\Gamma_{F_2}(x, y) = \frac{(1+x)(1+y)}{1-x-y-3xy} = \frac{G(x, y)}{H(x, y)}$$

and that the singularities of $\Gamma_{F_2}(x, y)$ are the points

$$\mathbf{z}' = (x, y) = \left(\frac{3p-2+2\sqrt{3p^2-3p+1}}{3p}, \frac{1-3p+2\sqrt{3p^2-3p+1}}{3(1-p)} \right)$$

whose coordinates are real numbers with positive coordinates. Hence,

$$(\mathbf{z}')^{-n \cdot \mathbf{r}} = \left(\frac{2-3p+2\sqrt{3p^2-3p+1}}{p}, \frac{3p-1+2\sqrt{3p^2-3p+1}}{(1-p)} \right)^{n \cdot \mathbf{r}} = e^{n \cdot \psi_{F_2}(\mathbf{r})}$$

We quickly check that the assumption of Theorem 5.1 from [16] are satisfied. In other words, we need to check that the partial derivative $\frac{\partial H}{\partial y}$ does not vanish at \mathbf{z}' and that the matrix \mathcal{H} from Equation (5.25) in [16] is non singular (with $\mathbf{w} = \mathbf{z}'$). Indeed, a direct computation gives

$$\frac{\partial H}{\partial y}(x, y) = -1 - 3x$$

which is non zero at \mathbf{z}' . Now, the dimension d being two, still with the notation of [16], the matrix \mathcal{H} is the scalar

$$\begin{aligned}\mathcal{H} &= V_1 + V_1^2 + U_{1,1} - 2V_1U_{1,2} + V_1^2U_{2,2} \\ &= \frac{x(1+3y)}{y(1+3x)} + \left(\frac{x(1+3y)}{y(1+3x)}\right)^2 + 0 - 2\frac{x(1+3y)}{y(1+3x)}\frac{3xy}{y(1+3x)} + 0 \\ &= \frac{xy + 3x^2y + 3xy^2 + x^2}{y^2(1+3x)^2} > 0.\end{aligned}$$

Therefore, following ([16], Equation 5.1) we get

$$f_{n,\mathbf{r}} = ce^{n\cdot\psi_{F_2}(\mathbf{r})}n^{-\frac{1}{2}}\left(1 + O\left(\frac{1}{n}\right)\right) \quad (43)$$

where $c = c(p)$ does not depend on p and \mathbf{r} is rational. In fact the results of [9] allows to have relation (43) when \mathbf{r} is irrational, only the left hand side should be replaced by the sum of the coefficients γ_i in the uniformly bounded neighborhood of the point $n\mathbf{r}$.

9 Concluding remarks and open questions

Finding of $\psi_{F_m}(\mathbf{r})$, where $\mathbf{r} = (p, q, 1 - p - q)$ for free group F_3 of rank 3 leads to the solving of polynomial equation of degree 4 in variable $z = e^s$

$$\begin{aligned}3p^2z^4 + 4p(7p - 2)z^3 + 2(33p^2 - 32pq - 8p - 32q^2 + 32q - 8)z^2 \\ + 12p(5p - 6)z - 45p^2 = 0\end{aligned} \quad (44)$$

and it can be solved in radicals. Substituting $p = q = \frac{1}{3}$ in (44) we obtain that $s = \log 5$ and hence we get a value $\psi_{F_3}(\frac{1}{3}, \frac{1}{3}, \frac{1}{3}) = \log 5$. So the multivariate growth in this case coincides with the ordinary growth and $\log 5$ is a maximal value of $\psi(\mathbf{r})$. The higher ranks $m = 4, 5, \dots$ lead to polynomial equations of degree > 5 and most probably obtaining of the precise analytic expressions for $\psi_{F_m}(\mathbf{r})$ is impossible. But at least we know that $\psi_{F_m}(\mathbf{r})$ is a real analytic concave function [9] with a maximum value $\log(2m - 1)$ achieved at unique point $\mathbf{r} = (\frac{1}{m}, \dots, \frac{1}{m})$.

Now, let us go back to cogrowth. It can be shown that the condition (CG) always hold for a subgroup $H < F_m$ and so the formula (4) is applicable. If $H < F_m$ is a finitely generated subgroup then it is represented by a regular language [4] and hence, its cogrowth and multivariate cogrowth series are rational. The Conjecture claiming that $\Gamma_H(z), z \in \mathbb{C}$ is rational if and only if H is finitely generated was stated in [4] and it is known that this conjecture is true in the case of normal subgroups. A similar conjecture can be stated for multivariate cogrowth series $\Gamma_H(\mathbf{z}), \mathbf{z} \in \mathbb{C}^m$. Also, we state the following:

Conjecture 9.1. *Let $N \triangleleft F_m$. Then F_m/N is amenable if and only if $\psi_N(\mathbf{r}) = \psi_{F_m}(\mathbf{r})$.*

By cogrowth criteria of amenability we know that in the case when F_m/N is amenable, the relation

$$\log(2m - 1) = \max_{\mathbf{r}} \psi_{F_m}(\mathbf{r}) = \max_{\mathbf{r}} \psi_N(\mathbf{r}) \quad (45)$$

hold. It is unclear if $\psi_N(\mathbf{r})$ may have values less than the values of $\psi_{F_m}(\mathbf{r})$ in the case when the Equation (45) holds. Even the case when $N = [F_2, F_2]$ is a commutator subgroup of F_2 deserves a separate consideration.

And finally, there is a formula in [1]

$$\chi(p) = 2 \min_t \left[\sum_{i=1}^m \sqrt{t^2 + p_i^2} - (m-1)t \right]$$

for the spectral radius $\chi(p)$ of a symmetric random walk on a free group F_m given by a positive vector $p = (p_1, \dots, p_m)$, $2 \sum p_i = 1$ where $p(a_i) = p(a_i^{-1}) = p_i$.

Computation of $\chi(p)$ in the case of rank 2 leads to the equation of degree 4 in variable $x = t^2$

$$3x^4 + 4(p_1^2 + p_2^2)x^3 + 6p_1^2p_2^2x^2 - p_1^4p_2^4 = 0 \quad (46)$$

and hence $\chi(p)$ can be expressed in radicals. Taking $p_1 = p_2 = \frac{1}{4}$, (46) leads to the equation

$$(1 + 16x)^3(-1 + 48x) = 0$$

which gives a value $\chi(p) = \frac{\sqrt{3}}{2}$.

The later number is known since 1959 due to H. Kesten [14] who, in particular proved that for a simple random walk on F_m the spectral radius $\chi = \frac{\sqrt{2m-1}}{m}$. Higher rank leads to solving polynomial equations of degree > 5 and expressing $\chi(p)$ in radicals seems to be impossible for $F_m, m \geq 3$ and arbitrary p .

Let $H < F_m$ and $\chi_{F_m/H}(p)$ be a spectral radius of a random walk on a Schreier graph $\Lambda = \Lambda(F_m, H, \Sigma)$ given by probabilities $p_i, 1 \leq i \leq m$. We end up with the following question.

Problem 9.1. *Is there a formula expressing $\chi_{F_m/H}$ via $\alpha_H(p)$, where $\alpha_H(p) = \varphi_H(p)$ is a multivariate growth of $\Delta_H(2p)$ in the direction prescribed by the vector $2p \in M_m$? Does such a formula exists when H is normal subgroup in F_m and hence $\Lambda = \Lambda(F_m, H, \Sigma)$ is a Cayley graph.*

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